

Challenges in the Application of Mathematical Programming in the Enterprise-wide Optimization of Process Industries

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Motivation for Enterprise-wide Optimization

Process industry:

Trend to global operations

Pressure for reducing costs, inventories and ecological footprint



Major goal: Enterprise-wide Optimization

Recent research area in Process Systems Engineering:

Grossmann (2005); Varma, Reklaitis, Blau, Pekny (2007)

A major challenge: optimization models and solution methods

EWO involves optimizing the operations of R&D, material supply, manufacturing, distribution of a company to reduce costs, inventories, ecological footprint and to maximize profits, responsiveness.

Key element: Supply Chain

Example: petroleum industry



Wellhead



Trading



Transfer of
Crude



Refinery
Processing



Schedule
Products



Transfer of
Products



Terminal
Loading

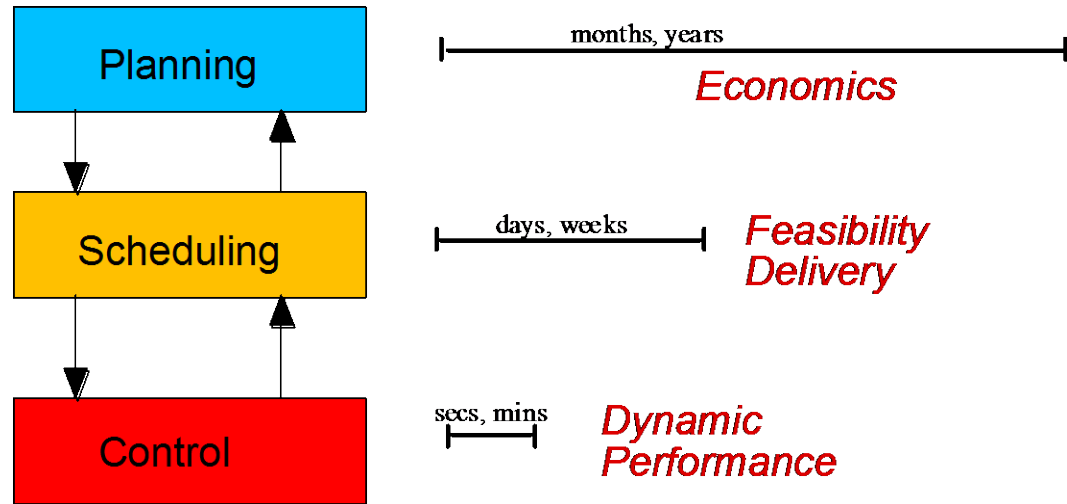


Pump

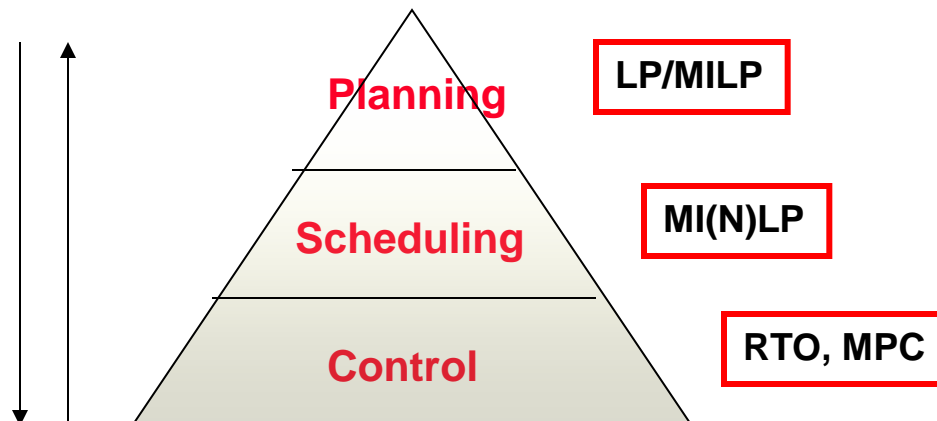
Key issues:

I. Integration of planning, scheduling and control

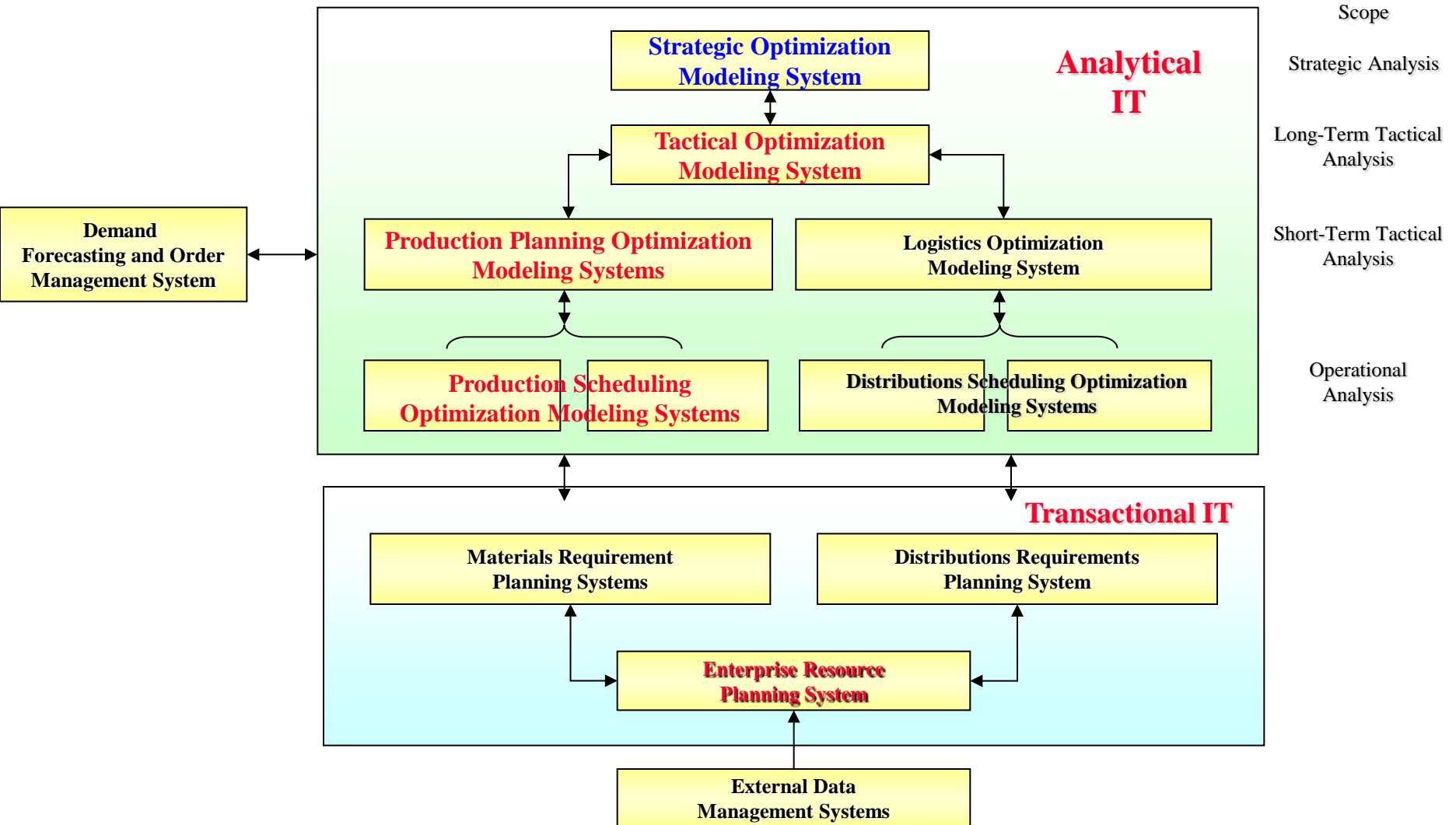
Multiple time scales



Multiple models



II. Integration of information and models/solution methods



Optimization Modeling Framework: Mathematical Programming

$$\min Z = f(x, y) \quad \text{Objective function}$$

$$s.t. \quad h(x, y) = 0 \quad \text{Constraints}$$

$$g(x, y) \leq 0$$

$$x \in R^n, \quad y \in \{0,1\}^m$$

MINLP: Mixed-integer Nonlinear Programming Problem

$$\begin{aligned} \min Z &= f(x) \\ \text{s.t. } h(x) &= 0 \\ g(x) &\leq 0 \\ x &\in R^n \end{aligned}$$

LP Codes:

CPLEX, XPRESS, GUROBI, XA

Very large-scale models

Interior-point: solvable polynomial time

NLP Codes:

CONOPT *Drud (1998)*

IPOPT *Wächter & Biegler (2006)*

Knitro *Byrd, Nocedal, Waltz (2006)*

MINOS *Murtagh, Saunders (1995)*

SNOPT *Gill, Murray, Saunders (2006)*

BARON *Sahinidis et al. (1998)*

Couenne *Belotti, Margot (2008)*

} Global
Optimization

Large-scale models

RTO: *Marlin, Hrymak (1996)*

Zavala, Biegler (2009)

Issues:

Convergence

Nonconvexities

$$\min Z = f(x, y)$$

$$s.t. \quad h(x, y) = 0$$

$$g(x, y) \leq 0$$

$$x \in R^n, \quad y \in \{0,1\}^m$$

MILP Codes:

CPLEX, XPRESS, GUROBI, XA

*Great Progress over last decade despite NP-hard
Planning/Scheduling: Lin, Floudas (2004)
Mendez, Cerdá, Grossmann, Harjunkski (2006)
Pochet, Wolsey (2006)*

MINLP Codes:

DICOPT (GAMS) Duran and Grossmann (1986)

a-ECP Westerlund and Petersson (1996)

MINOPT Schweiger and Floudas (1998)

MINLP-BB (AMPL) Fletcher and Leyffer (1999)

SBB (GAMS) Bussieck (2000)

Bonmin (COIN-OR) Bonami et al (2006)

FilMINT Linderoth and Leyffer (2006)

BARON Sahinidis et al. (1998)

Couenne Belotti, Margot (2008)

GLOMIQO Floudas, Meisner (2011)

Global
Optimization

*New codes over last decade
Leveraging progress in MILP/NLP*

Issues:

Convergence

Nonconvexities

Scalability

Modeling systems

Mathematical Programming

GAMS (*Meeraus et al, 1997*)

AMPL (*Fourer et al., 1995*)

AIMSS (*Bisschop et al. 2000*)

1. Algebraic modeling systems => pure equation models
2. Indexing capability => large-scale problems
3. Automatic differentiation => no derivatives by user
4. Automatic interface with
LP/MILP/NLP/MINLP solvers

Have greatly facilitated development and implementation of Math Programming models

Generalized Disjunctive Programming (GDP)

$$\min Z = \sum_k c_k + f(x)$$

Raman, Grossmann (1994)

$$s.t. \quad r(x) \leq 0$$

$$\bigvee_{j \in J_k} \begin{bmatrix} Y_{jk} \\ g_{jk}(x) \leq 0 \\ c_k = \gamma_{jk} \end{bmatrix} \quad k \in K$$

Disjunctions

$$\Omega(Y) = true$$

Logic Propositions

$$x \in R^n, \quad c_k \in R^1$$

Continuous Variables

$$Y_{jk} \in \{ true, false \}$$

Boolean Variables

Codes:

LOGMIP (*GAMS-Vecchiotti, Grossmann, 2005*)

EMP (*GAMS-Ferris, Meeraus, 2010*)

Framework for deriving scheduling models

Castro, Grossmann (2012)

Optimization Under Uncertainty

Multistage Stochastic Programming

Birge & Louveaux, 1997; Sahinidis, 2004

$$\min z = c^1 x^1 + E_{\omega^2} [c^2(\omega) x^2(\omega^2) + \dots + E_{\omega^N} [c^N(\omega) x^N(\omega^N)] \dots]$$

$$\text{s.t.} \quad W^1 x^1 = h^1$$

$$T^1(\omega) x^1 + W^2 x^2(\omega^2) = h^2(\omega)$$

$$\vdots$$

*Exogeneous uncertainties
(e.g. demands)*

$$T^{N-1}(\omega) x^{N-1}(\omega^{N-1}) + W^N x^N(\omega^N) = h^N(\omega)$$

$$x^1 \geq 0, x^t(\omega^t) \geq 0, t = 2, \dots, N-1$$

Special case: two-stage programming (N=2)

$$\boxed{x^1 \text{ stage 1} \quad \omega \quad x^2 \text{ recourse (stage 2)}}$$

Planning with endogenous uncertainties (e.g. yields, size reservoir, test drug):

Goel, Grossmann (2006), Colvin, Maravelias (2009), Gupta, Grossmann (2011)

Robust Optimization

Major concern: **feasibility over uncertainty set**

Ben-Tal et al., 2009; Bertsimas and Sim (2003)

LP

$$\min_x c^T x \quad : \quad a_i^T x \leq b_i, \quad i = 1, \dots, m$$

- Ellipsoidal uncertainty:

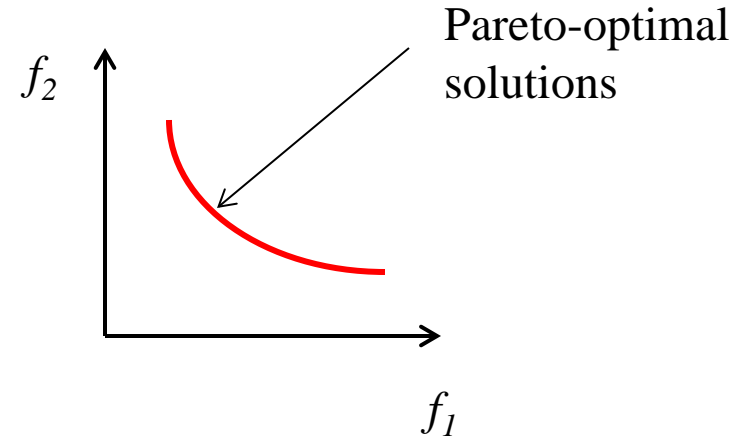
$$a_i \in \mathcal{E}_i = \{\hat{a}_i + P_i^{1/2} u \quad : \quad \|u\|_2 \leq 1\}$$

Robust scheduling:

Lin, Janak, Floudas (2004); Li, Ierapetritou (2008)

Multiobjective Optimization

$$\begin{aligned} \min Z &= \begin{bmatrix} f_1(x, y) \\ f_2(x, y) \\ \dots\dots \end{bmatrix} \\ \text{s.t. } & h(x, y) = 0 \\ & g(x, y) \leq 0 \\ & x \in R^n, \quad y \in \{0,1\}^m \end{aligned}$$



ϵ -constraint method: Ehrgott (2000)

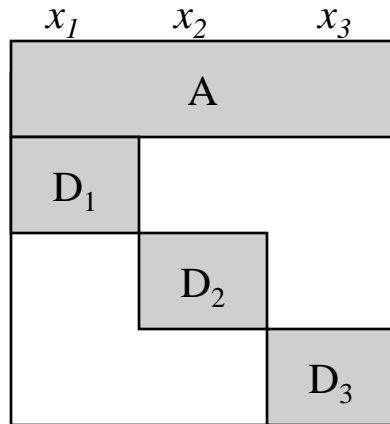
Parametric programming: Pistikopoulos, Georgiadis and Dua (2007)

Decomposition Techniques

Lagrangian decomposition

Geoffrion (1972) Guinard (2003)

Complicating Constraints



complicating constraints \rightarrow

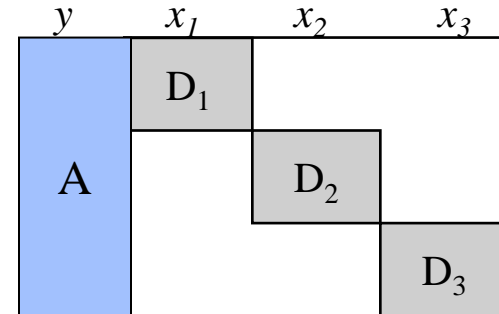
$$\begin{aligned} \max \quad & c^T x \\ \text{st} \quad & Ax = b \\ & D_i x_i = d_i, \quad i = 1, \dots, n \\ & x \in X = \{x \mid x_i, i = 1, \dots, n, |x_i \geq 0\} \end{aligned}$$

Widely used in EWO

Benders decomposition

Benders (1962), Magnanti, Wing (1984)

Complicating Variables



complicating variables \rightarrow

$$\begin{aligned} \max \quad & a^T y + \sum_{i=1, \dots, n} c_i^T x_i \\ \text{st} \quad & Ay + D_i x_i = d_i, \quad i = 1, \dots, n \\ & y \geq 0, x_i \geq 0, \quad i = 1, \dots, n \end{aligned}$$

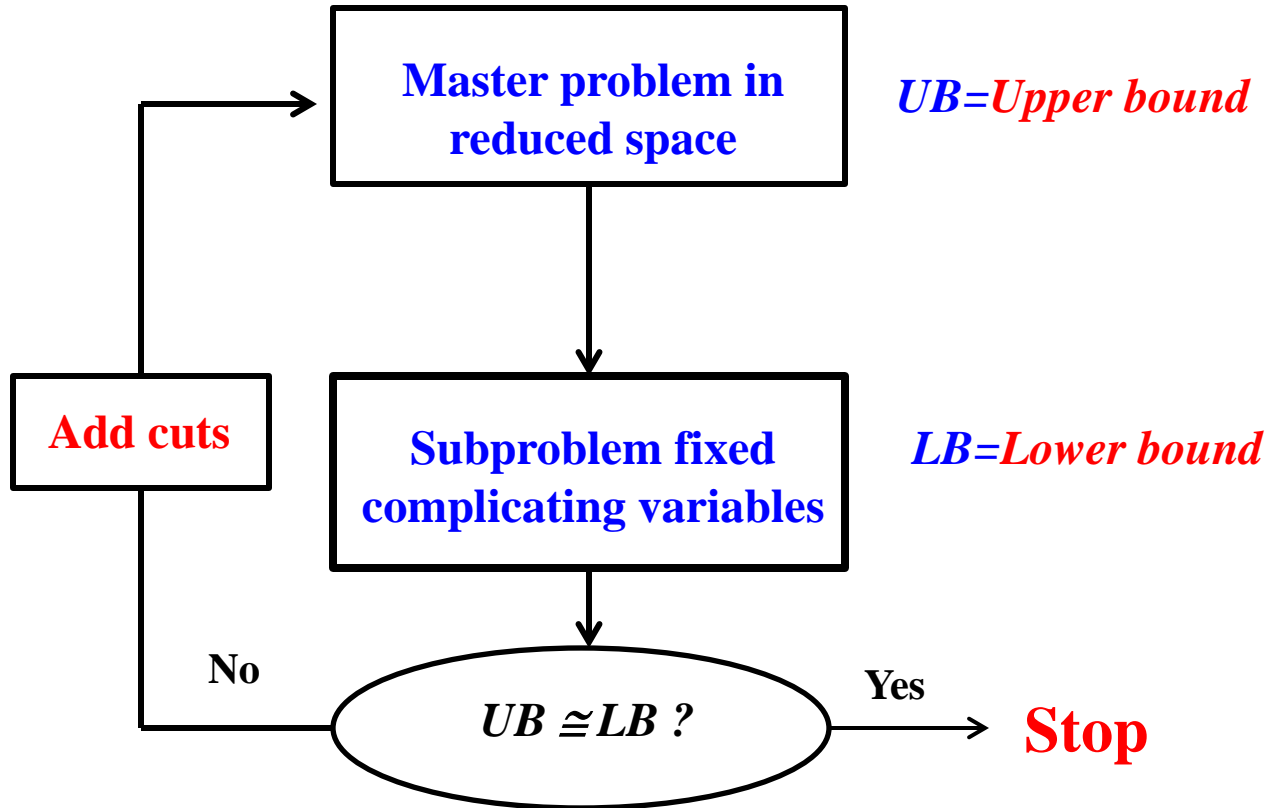
Applied in 2-stage Stochastic Programming

Decomposition Techniques (*cont.*)

Bi-level decomposition

Tailor-made Benders

Iyer, Grossmann (1998)



**Special industrial interest group in CAPD:
“Enterprise-wide Optimization for Process Industries”**

<http://egon.cheme.cmu.edu/ewocp/>

**Multidisciplinary team:
Chemical engineers, Operations Research, Industrial Engineering**

Researchers:

Carnegie Mellon:

- Ignacio Grossmann (ChE)
- Larry Biegler (ChE)
- Nicola Secomandi (OR)
- John Hooker (OR)

Carnegie Mellon



Lehigh University:

- Katya Scheinberg (Ind. Eng)
- Larry Snyder (Ind. Eng.)
- Jeff Linderoth (Ind. Eng.)



Carnegie Mellon

Projects and case studies with partner companies: “Enterprise-wide Optimization for Process Industries”

ABB: *Optimal Design of Supply Chain for Electric Motors*

Contact: Iiro Harjunkoski

Ignacio Grossmann, Analia Rodriguez, Yonheng Jiang

Air Liquide: *Optimal Coordination of Production and Distribution of Industrial Gases*

Contact: Jean Andre, Jeffrey Arbogast

Ignacio Grossmann, Pablo Marchetti

Braskem: *Optimal production and scheduling of polymer production*

Contact: Rita Majewski, Wiley Bucey

Ignacio Grossmann, Pablo Marchetti

Dow: *Optimal Design of Supply Chains under Disruptions*

Contact: John Wassick

Ignacio Grossmann, Pablo Garcia-Guerrero

Dow: *Optimal Operation of Reliable Integrated Sites*

Contact: John Wassick, Anshul Agrawal

Ignacio Grossmann, Bruno Calfa

Dow: *Financial Risk with Discrete Event Simulation*

Contact: Bikram Shards, Scott Bury

Nikolaos Sahinidis, Sayit Amaran

Dow: *Batch Scheduling and Dynamic Optimization*

Contact: Carlos Villa

Larry Biegler, Yisu Nie

Ecopetrol: *Adaptive Process Control*

Contact: Sandra Milena Montagut

Erik Ydstie, Masoud Golshan

ExxonMobil: *Global optimization of multiperiod blending networks*

Contact: Myun-Seok Cheon, Kevin Furman, Nick Sawaya

Ignacio Grossmann, Scott Kolodziej, Francisco Trespalacios

ExxonMobil: *Design and planning of oil and gasfields with fiscal constraints*

Contact: Bora Tarhan

Ignacio Grossmann, Vijay Gupta

Mitsubishi: *Optimization of power flows*

Contact: Arvind Raghunathan

Larry Biegler, Ajit Gopalakrishnan

Petrobras: *Nonlinear Integrated Model for Operational Planning of Multi-Site Refineries*

Contact: Lincoln Moro

Ignacio Grossmann, Breno Menezes

P&G: *Models for predicting shelf-life of consumer products*

Contact: Ben Weinstein

Larry Biegler, George Ostace

Praxair: *Capacity Planning of Power Intensive Networks with Changing Electricity Prices*

Contact: Jose Pinto

Ignacio Grossmann, Sumit Mitra

UNILEVER: *Planning and Scheduling of Fast Moving Goods*

Contact: Hans Hogland

Ignacio Grossmann, Martijn van Elzakker

Major Issues

- **Linear vs Nonlinear models**
- **The multi-scale optimization challenge**
- **The uncertainty challenge**
- **Economics vs performance**
- **Computational efficiency in large-scale problems**
- **Commercial vs. Tailored Software**

-Linear vs Nonlinear Models

Most EWO problems formulated as MILP

Example: MILP Supply Chain Design Problem

2,001 0-1 vars, 37,312 cont vars, 80,699 constraints

CPLEX 12.2:

MIP Solution: 5,043,251 (160 nodes, 13734 iterations,)

Relative gap: 0.004263 (< 0.5%)

CPU-time: 27 secs!!!

NLP *required for process models*

MINLP *required for cyclic scheduling, stochastic inventory, MIDO for integration of control*

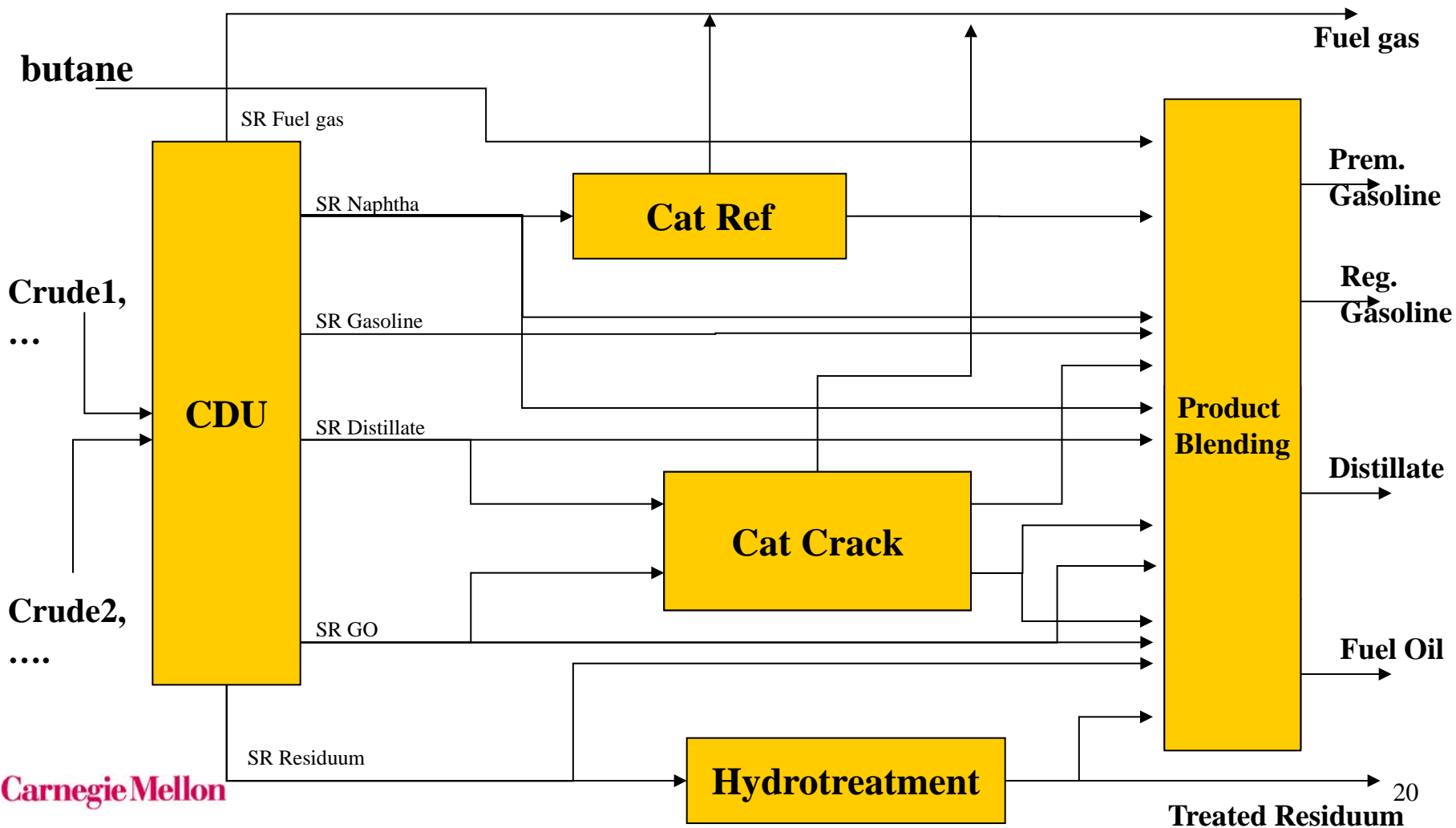
Nonlinear CDU Models in Refinery Planning Optimization



Alattas, Palou-Rivera, Grossmann (2010)

Typical Refinery Configuration

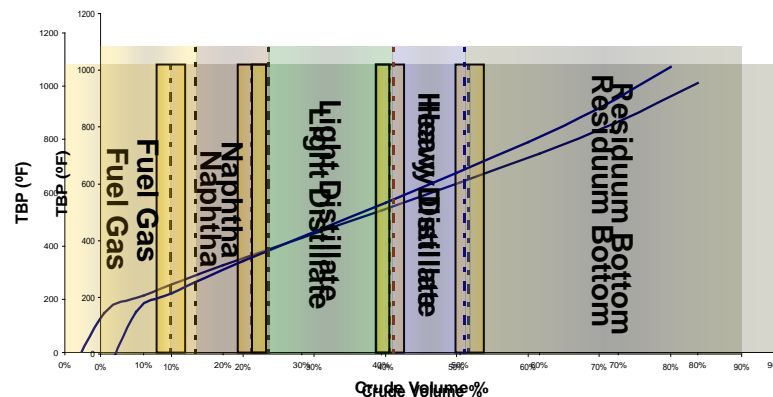
(Adapted from Aronofsky, 1978)



Refinery Planning Models

LP planning models

Fixed yield model
Swing cuts model



Nonlinear FI Model (*Fractionating Index*)

- FI Model is crude independent
 - *FI values are characteristic of the column*
 - *FI values are readily calculated and updated from refinery data*
- Avoids more complex, nonlinear modeling equations
- Generates cut point temperature settings for the CDU
- Adds few additional equations to the planning model

Planning Model Example Results

Crude1	Louisiana	Sweet	Lightest
Crude2	Texas	Sweet	↓
Crude3	Louisiana	Sour	
Crude4	Texas	Sour	Heaviest

- Comparison of *nonlinear fractionation index (FI)* with the fixed yield (FY) and swing cut (SC) models
- Economics: maximum profit

FI yields highest profit

Model	Case1	Case2	Case3
FI	245	249	247
SC	195	195	191
FY	51	62	59



Model statistics LP vs NLP

- FI model larger number of equations and variables
- Impact on solution time
- ~30% nonlinear variables

	Model	Variables	Equations	Nonlinear Variables	CPU Time	Solver
2 Crude Oil Case	<i>FY</i>	128	143		0.141	CPLEX
	<i>SC</i>	138	163		0.188	
	<i>FI</i>	1202	1225	348	0.328	CONOPT

- **Solution large-scale problems:**

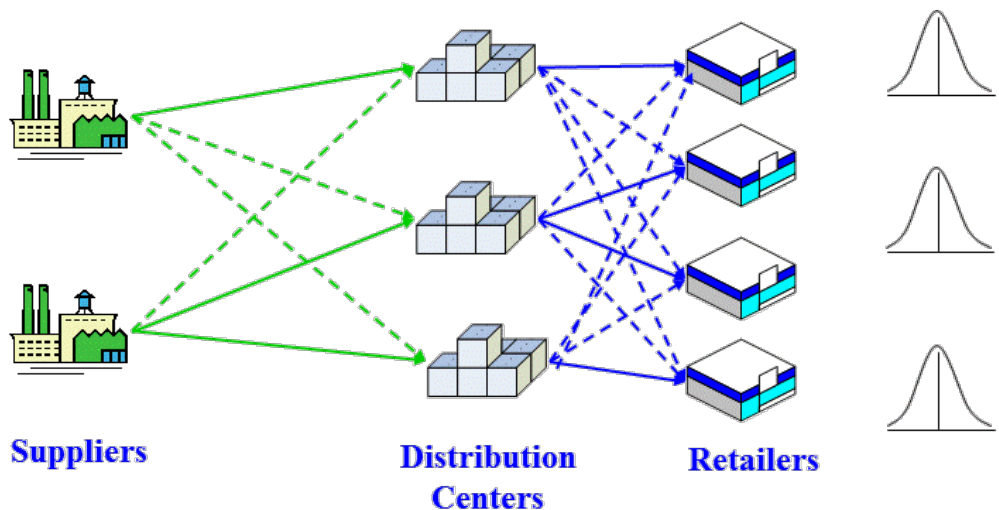
Strategy 1: Exploit problem structure (TSP)

Strategy 2: Decomposition

*Strategy 3: Heuristic methods to obtain
“good feasible solutions”*

Design Supply Chain Stochastic Inventory

You, Grossmann (2008)



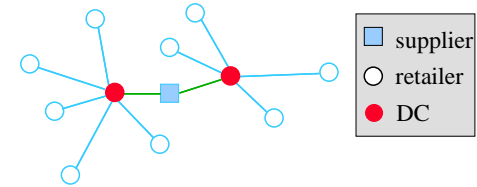
- Major Decisions (Network + Inventory)
 - ◆ Network: number of DCs and their locations, assignments between retailers and DCs (single sourcing), shipping amounts
 - ◆ Inventory: number of replenishment, reorder point, order quantity, safety stock

- Objective: (Minimize Cost)
 - ◆ Total cost = DC installation cost + transportation cost + fixed order cost + working inventory cost + safety stock cost

Trade-off: Transportation vs inventory costs

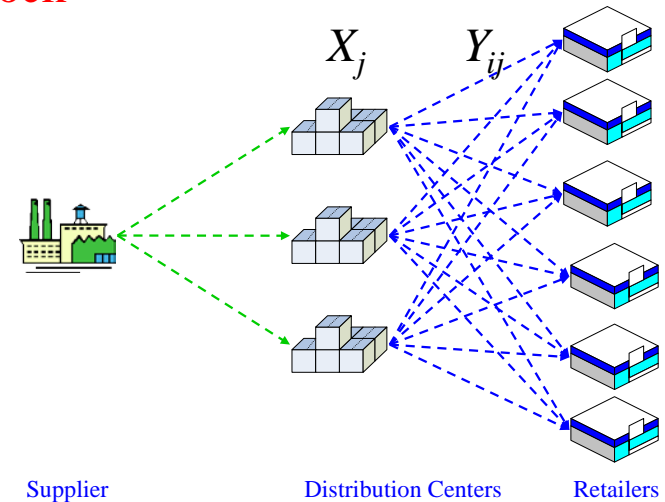
INLP Model Formulation

$$\begin{aligned}
 \min \quad & \sum_{j \in J} f_j X_j && \text{DC installation cost} \\
 & + \beta \sum_{j \in J} \sum_{i \in I} d_{ij} \chi \mu_i Y_{ij} && \text{DC - retailer transportation} \\
 & + \sum_{j \in J} \sqrt{2\theta h (F_j + \beta g_j) \sum_{i \in I} \chi \mu_i Y_{ij}} + \beta \sum_{i \in I} \sum_{j \in J} (a_j \chi \mu_i Y_{ij}) && \text{EOQ} \\
 & + \sum_{j \in J} (\theta h z_\alpha \sqrt{\sum_{i \in I} \sigma_i^2 L_i Y_{ij}}) && \text{Safety Stock}
 \end{aligned}$$



$$\begin{aligned}
 \text{s.t.} \quad & \sum_{j \in J} Y_{ij} = 1, \quad \forall i \in I \\
 & Y_{ij} \leq X_j, \quad \forall i \in I, j \in J \\
 & X_j, Y_{ij} \in \{0, 1\}, \quad \forall i \in I, j \in J
 \end{aligned}$$

} Assignments



Nonconvex INLP: 1. Variables Y_{ij} can be relaxed as continuous

2. Problem reformulated as MINLP

3. Solved by Lagrangean Decomposition (by distribution centers)

- Variables Y_{ij} can be relaxed as continuous variables (MINLP)
 - ◆ Local or global optimal solution always have all Y_{ij} at integer
 - ◆ If $h=0$, it reduces to an “uncapacitated facility location” problem
 - ◆ NLP relaxation is very effective (usually return integer solutions)

$$\begin{aligned}
 \min \quad & \sum_{j \in J} f_j X_j + \sum_{i \in I} \sum_{j \in J} \hat{d}_{ij} Y_{ij} + \sum_{j \in J} K_j \underbrace{\sqrt{\sum_{i \in I} \mu_i Y_{ij}}}_{Z1_j} + \sum_{j \in J} q \underbrace{\sqrt{\sum_{i \in I} \hat{\sigma}_i^2 Y_{ij}}}_{Z2_j} \\
 \text{s.t.} \quad & \sum_{j \in J} Y_{ij} = 1 \quad , \quad \forall i \in I \\
 & Y_{ij} \leq X_j \quad , \quad \forall i \in I, j \in J \\
 & Y_{ij} \geq 0 \quad , \quad \forall i \in I, j \in J \\
 & X_j \in \{0, 1\} \quad , \quad \forall j \in J
 \end{aligned}$$

Avoid unbounded gradient

Non-convex MINLP

where $\hat{d}_{ij} = \beta \mu_i (d_{ij} + a_j)$, $\hat{\sigma}_i^2 = L \sigma_i^2$, $K_j = \sqrt{2\theta h (F_j + \beta g_j)}$, $q = \theta h z_\alpha$

Algorithm 1 – MINLP Heuristic

- MINLP Heuristic Method
 - ◆ Solve the **convex relaxation** (MILP), **using secant for convex envelope**
 - ◆ Use optimal value of **X and Y** variables as **initial point**, solve the reformulated problem with an **MINLP solver** (BARON, Dicopt, etc.)

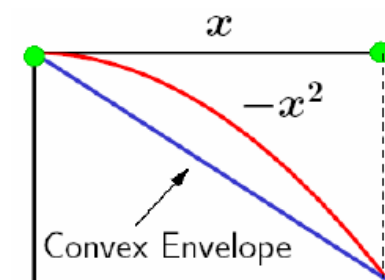
$$\min \sum_{j \in J} \left\{ f_j X_j + \left(\sum_{i \in I} \hat{d}_{ij} Y_{ij} \right) + K_j Z1_j + q Z2_j \right\}$$

$$\text{s.t.} \quad \sum_{j \in J} Y_{ij} = 1, \quad \forall i \in I$$

$$Y_{ij} \leq X_j, \quad \forall i \in I, j \in J$$

$$Y_{ij} \geq 0, \quad \forall i \in I, j \in J$$

$$X_j \in \{0, 1\}, \quad \forall j \in J$$



$$-Z1_j^2 + \sum_{i \in I} \mu_i Y_{ij} \leq 0, \quad \forall j \in J$$

$$-Z2_j^2 + \sum_{i \in I} \hat{\sigma}_i^2 Y_{ij} \leq 0, \quad \forall j \in J$$

$$Z1_j \geq 0, Z2_j \geq 0, \quad \forall j \in J$$

Convex
Relaxation



$$-\sqrt{\sum_{i \in I} \mu_i} \cdot Z1_j + \sum_{i \in I} \mu_i Y_{ij} \leq 0, \quad \forall j \in J$$

$$-\sqrt{\sum_{i \in I} \hat{\sigma}_i^2} \cdot Z2_j + \sum_{i \in I} \hat{\sigma}_i^2 Y_{ij} \leq 0, \quad \forall j \in J$$

Algorithm 2 - Lagrangean Relaxation

- Lagrangean Relaxation (LR) and Decomposition

- LR: dualizing the **single sourcing constraint**:

$$\sum_{j \in J} Y_{ij} = 1, \forall i \in I$$

- Spatial Decomposition**: decompose the problem for each potential DC j

- Implicit constraint**: at least one DC should be installed,

$$\sum_{j \in J} X_j \geq 1$$

- Use a special case of LR subproblem that $X_j=1$

$$\min \sum_{j \in J} \left\{ f_j X_j + \sum_{i \in I} (\hat{d}_{ij} - \lambda_i) Y_{ij} + K_j Z1_j + q Z2_j \right\} + \sum_{i \in I} \lambda_i$$

$$\text{s.t.} \quad Y_{ij} \leq X_j, \quad \forall j \in J, i \in I$$

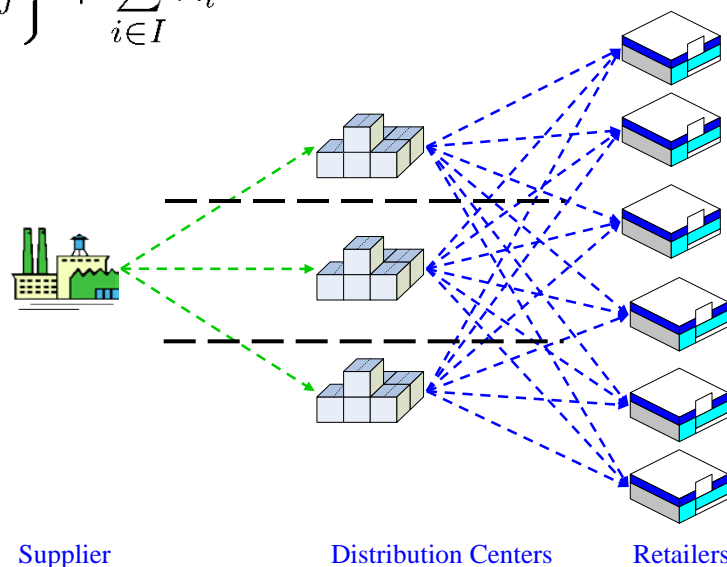
$$Y_{ij} \geq 0, \quad \forall j \in J, i \in I$$

$$X_j \in \{0, 1\}, \quad \forall j \in J$$

$$-Z1_j^2 + \sum_{i \in I} \mu_i Y_{ij} \leq 0, \quad \forall j \in J$$

$$-Z2_j^2 + \sum_{i \in I} \hat{\sigma}_i^2 Y_{ij} \leq 0, \quad \forall j \in J$$

$$Z1_j \geq 0, Z2_j \geq 0, \quad \forall j \in J$$



decompose by DC j

Computational Results

- Each instance has the same number of potential DCs as the retailers

150 retailers: MINLP has 150 bin. var., 22,800 cont. var., 22,800 constraints

No. Retailers	β	θ	Lagrangean Relaxation				
			Upper Bound	Lower Bound	Gap	Iter.	Time (s)
88	0.001	0.1	867.55	867.54	0.001 %	21	356.1
88	0.001	0.5	1230.99	1223.46	0.615 %	24	322.54
88	0.005	0.1	2284.06	2280.74	0.146 %	55	840.28
88	0.005	0.5	2918.3	2903.38	0.514 %	51	934.85
150	0.001	0.5	1847.93	1847.25	0.037 %	13	659.1
150	0.005	0.1	3689.71	3648.4	1.132 %	53	3061.2

- Suboptimal solution in 3 out of 6 cases with BARON for 10 hour limit. Large optimality gaps

The multi-scale optimization challenge

Temporal integration long-term, medium-term and short-term Bassett, Pekny, Reklaitis (1993), Gupta, Maranas (1999), Jackson, Grossmann (2003), Stefansson, Shah, Jenssen (2006), Erdirik-Dogan, Grossmann (2006), Maravelias, Sung (2009), Li and Ierapetritou (2009), Verderame, Floudas (2010), Salema, Barbosa-Povoa, Novais (2010)

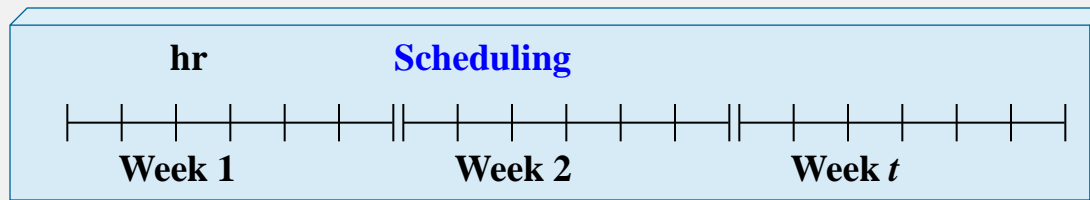
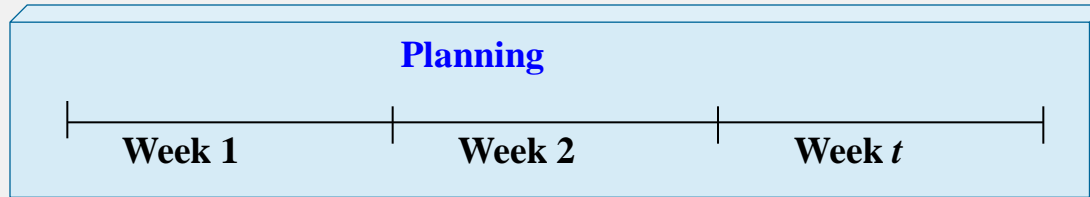
Spatial integration geographically distributed sites Gupta, Maranas (2000), Tsiakis, Shah, Pantelides (2001), Jackson, Grossmann (2003), Terrazas, Trotter, Grossmann (2011)

Decomposition is key: Benders, Lagrangean, bilevel

Multi-site planning and scheduling involves different temporal and spatial scales

Terrazas, Grossmann (2011)

Site 1



Weekly aggregate production:

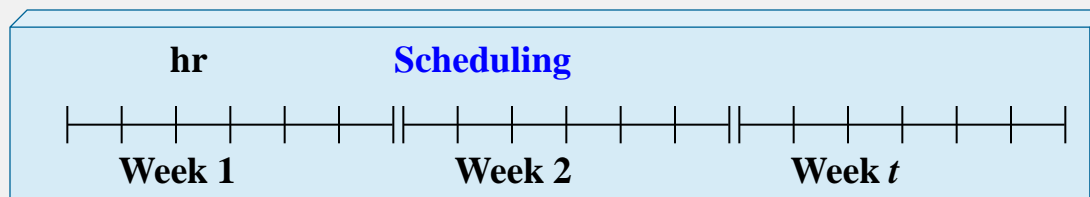
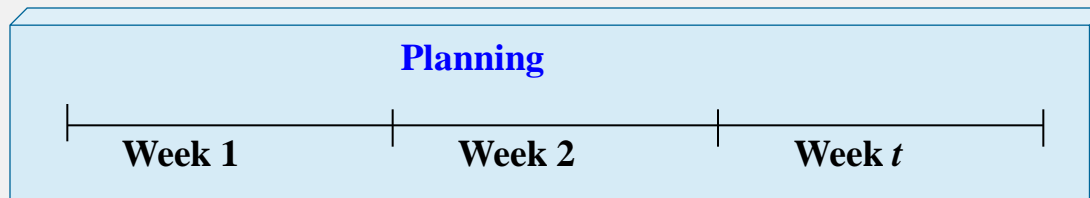
- Amounts
- Aggregate sequencing model (TSP constraints)

Detailed operation

- Start and end times
- Allocation to parallel lines

Different Temporal Scales

Site s



Weekly aggregate production:

Detailed operation

Different Spatial Scales

MILP Model

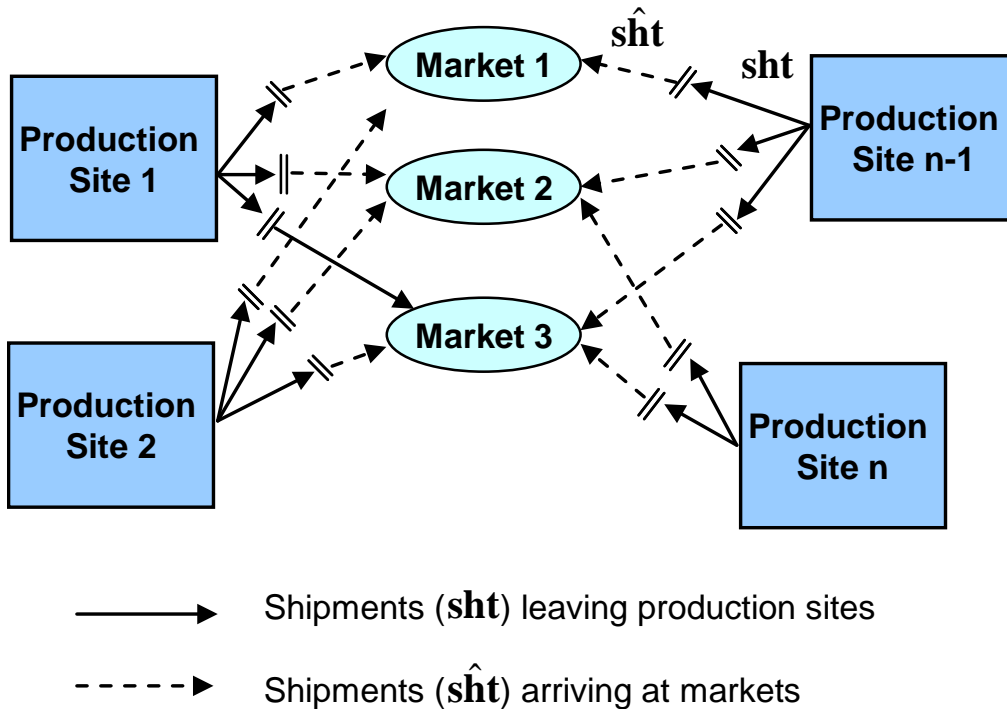
- **Objective: Maximize Profit**

subject to

- **Market constraints**
 - **Balance of sales vs. shipments to markets**
- **Production constraints**
 - **Capacity constraints:** Limited capacity at each production sites
 - **Inventory constraints:** Penalties for inventory over or under target
 - **Links across periods:** a) Carry over inventories from last month
b) Changeover to first product in next month
 - **Time Balances:** Task should not take longer than available time
 - **Sequencing Constraints:** Traveling Salesman Problem Constraints



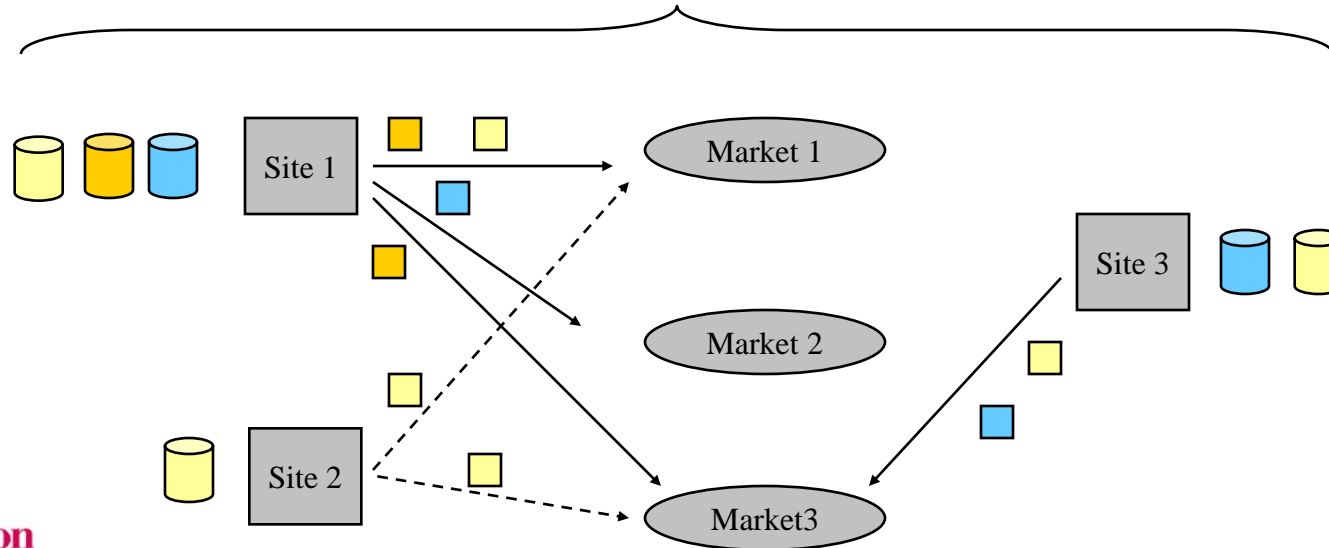
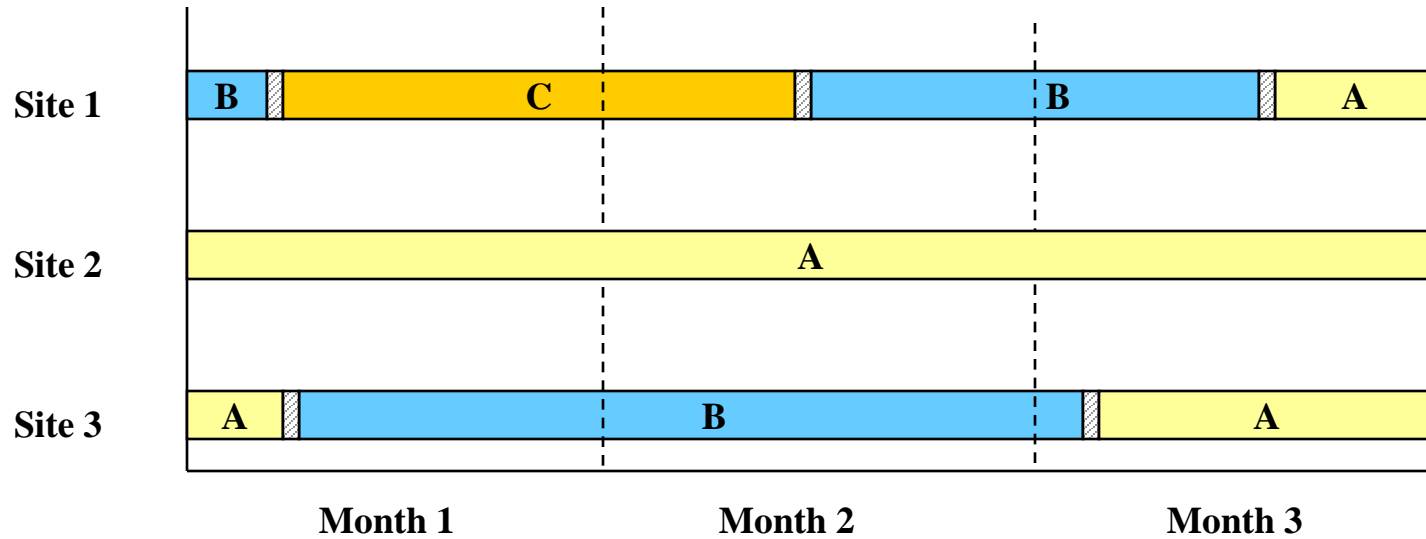
Source of complexity of the model: TSP constraints



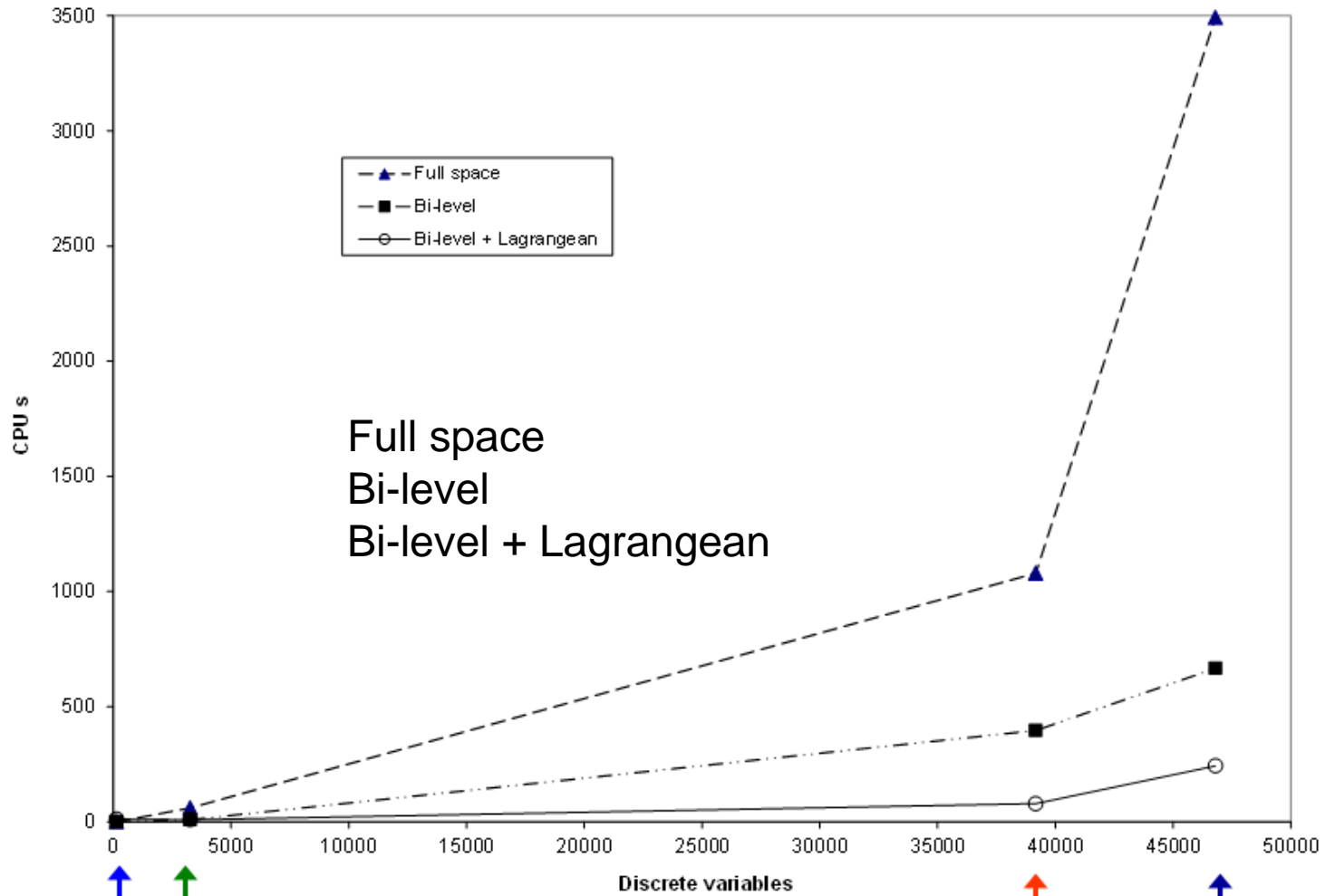
- **Bilevel decomposition**
 - Decouples **planning** from **scheduling**
 - Integrates across **temporal scale**
- **Lagrangean decomposition**
 - Decouples the solution of **each production site**
 - Integrates across **spatial scale**

Example: 3 sites, 3 products, 3 months

Profit: \$ 2.576 million



Large-scale problems



3 sites
2 markets
3 products
4 weeks

3 sites
2 markets
16 products
4 weeks

6 sites
6 markets
16 products
24 weeks

6 sites
6 markets
25 products
12 weeks

Mitra et al. (2013)



Given:

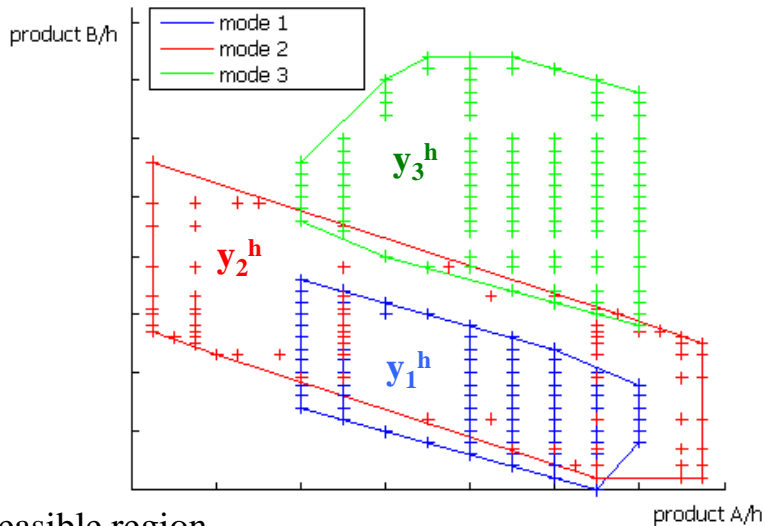
- Power-intensive plant
- Products $g \in G$ (Storable and Nonstorable)
- Product demands d_g^t for season $t \in T$
- Seasonal electricity prices on an hourly basis $e^{t,h}$, $t \in T$, $h \in H$
- Upgrade options $u \in U$ of existing equipment
- New equipment options $n \in N$
- Additional storage tanks $st \in ST$

Determine:

- Production levels $Pr_g^{t,h}$
 - Mode of operation $\tilde{y}_{m,o}^{t,h}, y_m^{t,h}$
 - Sales $S_g^{t,h}$
 - Inventory levels $INV_g^{t,h}$
- for each season on an hourly basis
- Upgrades for equipment $VU_{m,u}^t$
 - Purchase of new equipm. VN_n^t
 - Purchase of new tanks $VS_{st,g}^t$

With minimum investment and operating costs

Demand Side Management is part of a complex multi-scale design and operations problem for power-intensive processes.



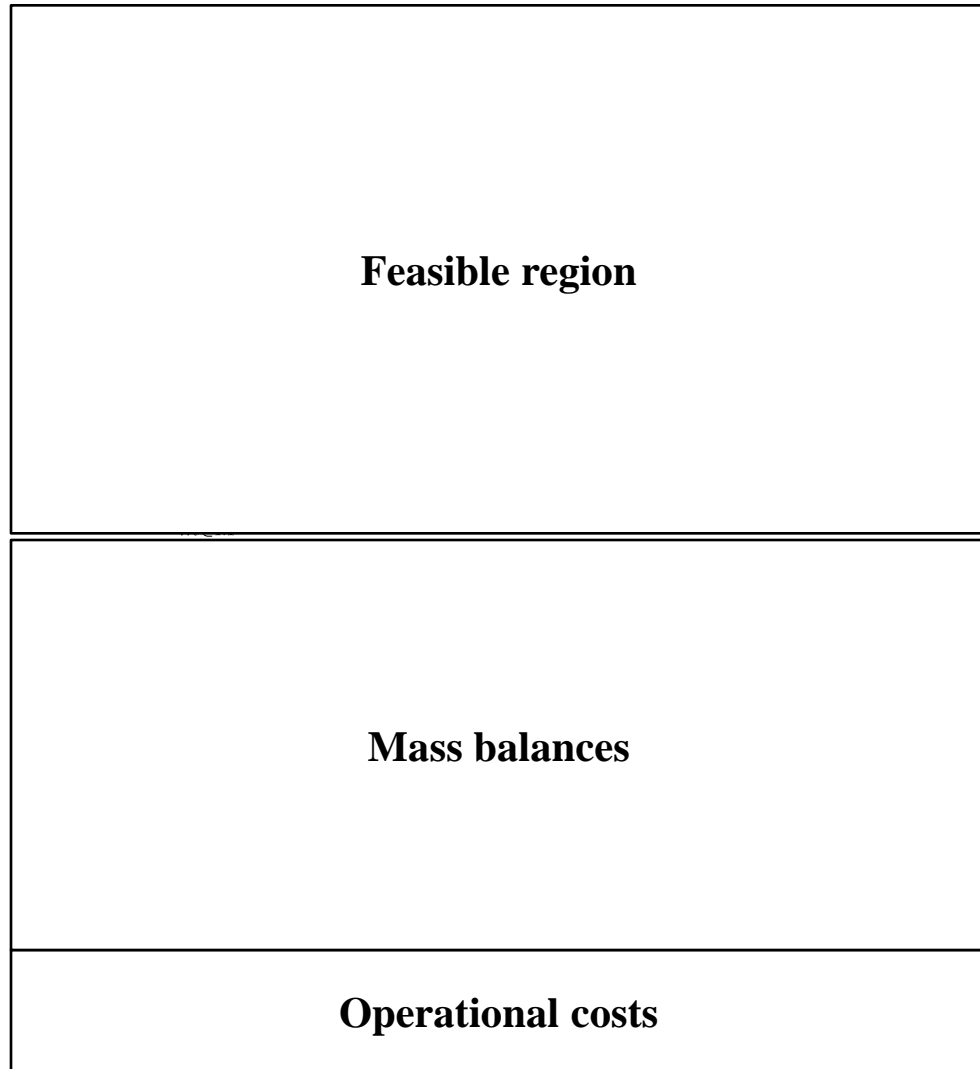
Feasible region

Feasible region: projection in product space

Modes: different ways of operating a plant

Mass balances: differences for products with and without inventory

Energy consumption: requires correlation with production levels for each mode

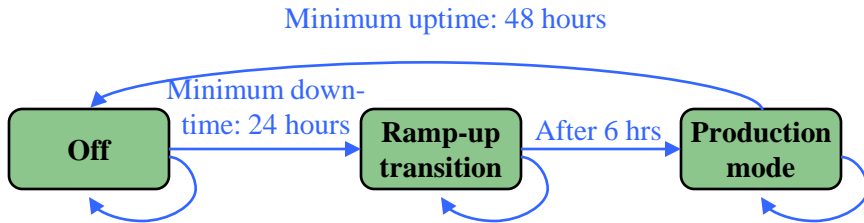


[1] Ierapetritou, M.G.; Wu, D.; Vin, J.; Sweeny P.; Chigirinskiy M. Cost Minimization in an Energy-Intensive Plant Using Mathematical Programming Approaches. Industrial & Engineering Chemistry Research, 41:5262–5277, 2002.

[2] Karwan, K.; Kebulis M. Operations planning with real time pricing of a primary input. Computers & Operations Research, 34:848–867, 2007.

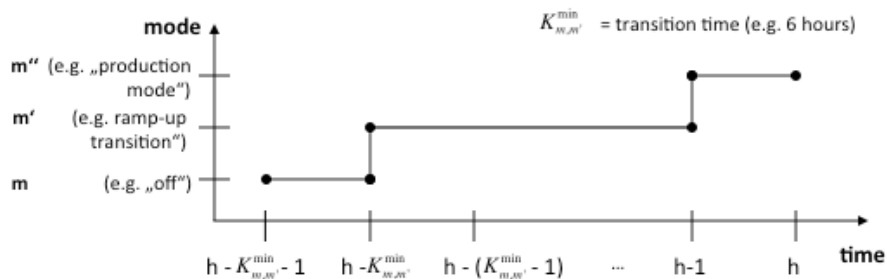
[3] Reformulation of disjunction with convex hull according to

Balas, E. Disjunctive Programming and a Hierarchy of Relaxations for Discrete Optimization Problems. SIAM J. Alg. Disc. Meth, 6:466–486, 1985



State diagram for transitions:

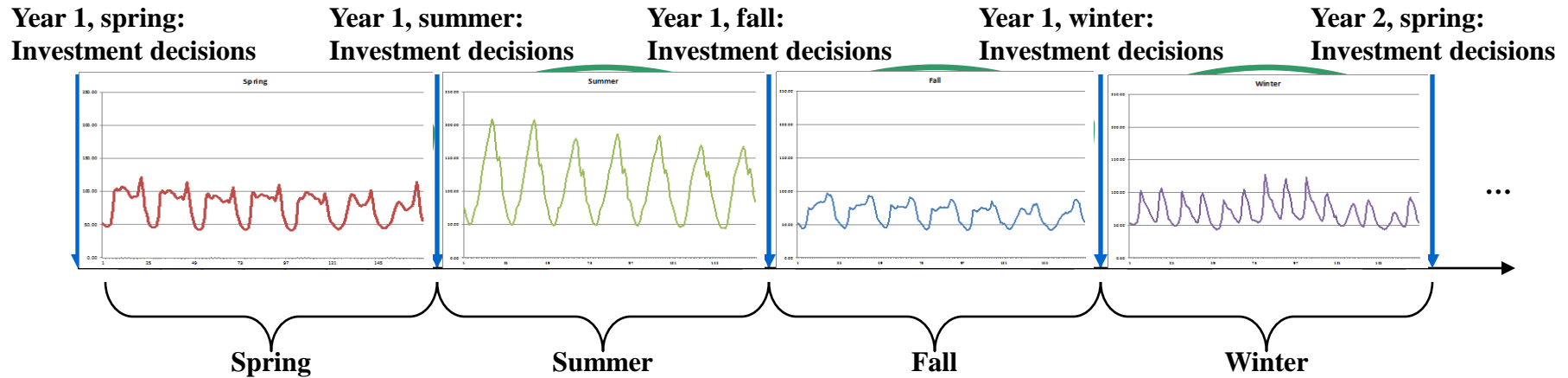
- nodes: states (**modes**) $y_{p,m}^h$
= different ways of operating a plant
- arcs = allowed **transitions** $z_{p,m,m'}^h$
(including constraints, e.g. min. up-/downtime)



**Logic constraints for transitions
between different modes**

Constraints for transitions within one mode

[*] Derivation of logic constraint using propositional logic according to Raman, R.; Grossmann, I.E. Modeling and Computational Techniques for Logic Based Integer Programming. Comp. Chem. Eng., 18:563, 1993.



- Horizon: 5-15 **years**, each year has 4 **periods** (spring, summer, fall, winter)
- Each period is represented by **one week on an hourly basis**
- Varying inputs: **electricity prices**, **demand data** (here: highly utilized plant), **configuration slates**
- Each representative week is repeated in a **cyclic** manner (**13** weeks reduced to **1** week)
(8736 hr vs. 672 hr)
- Connection between periods: Only through investment (design) decisions
- Design decisions are modeled by **discrete equipment sizes**

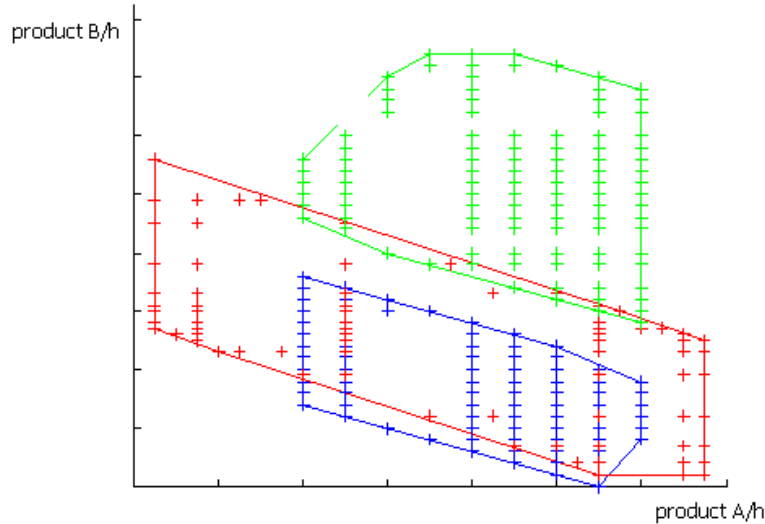


Fig. 1: Feasible region

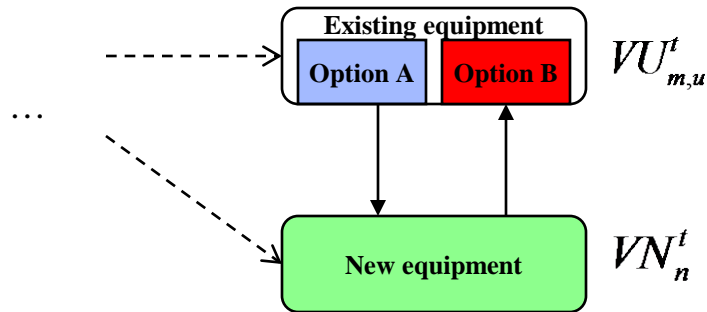


Fig. 2: State graph

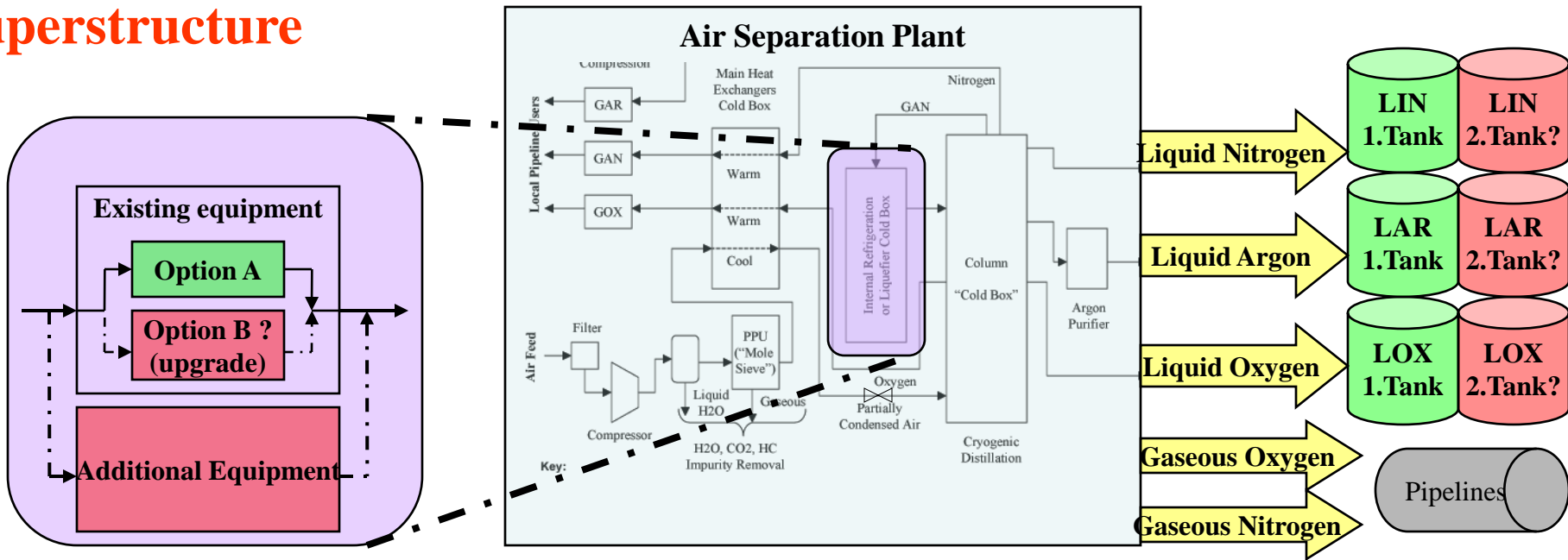
Logic constraints for equipment upgrades
(polyhedral representation of mode changes)

Logic constraints for additional equipment
(additional mode(s) are added to the state graph)

Constraints for additional storage
(manipulation of upper bound for inventory)

Objective: minimize investment + operational costs

Superstructure



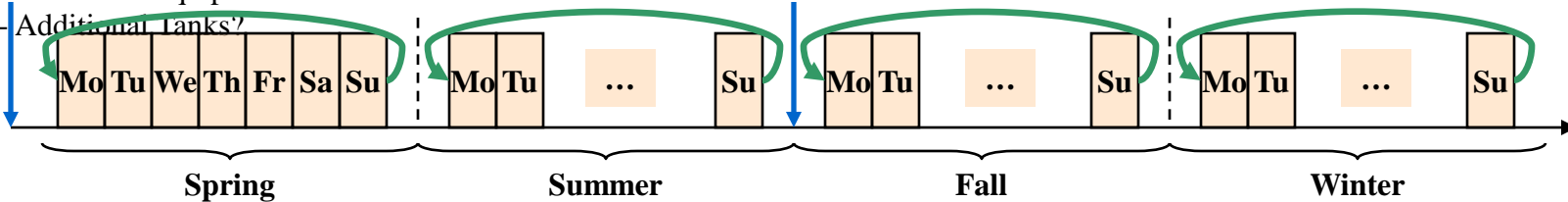
Time

Spring - Investment decisions:
(yes/no)

- Option B for existing equipment?
- Additional equipment?
- Additional Tanks?

Fall - Investment decisions: (yes/no)

- Option B for existing equipment?
- Additional equipment?
- Additional Tanks?



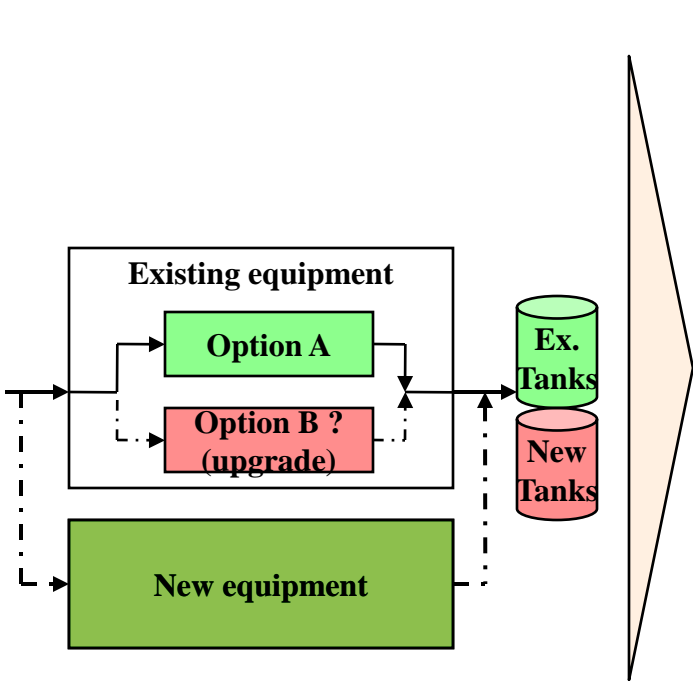


Fig. 1: Flowsheet superstructure

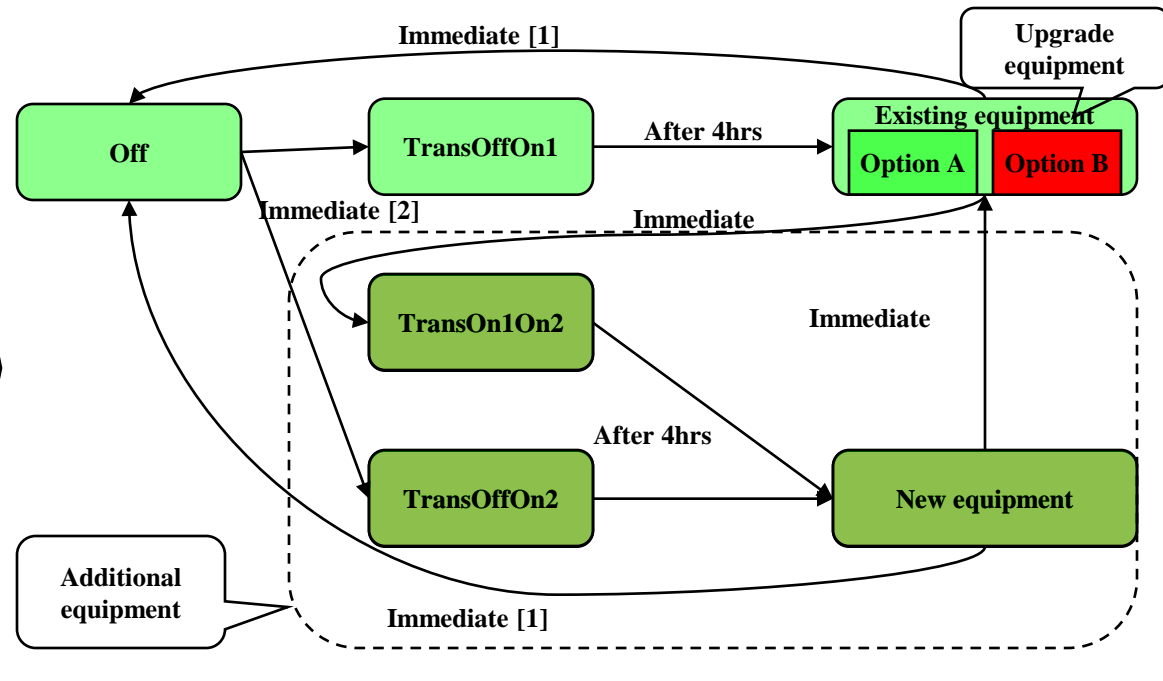
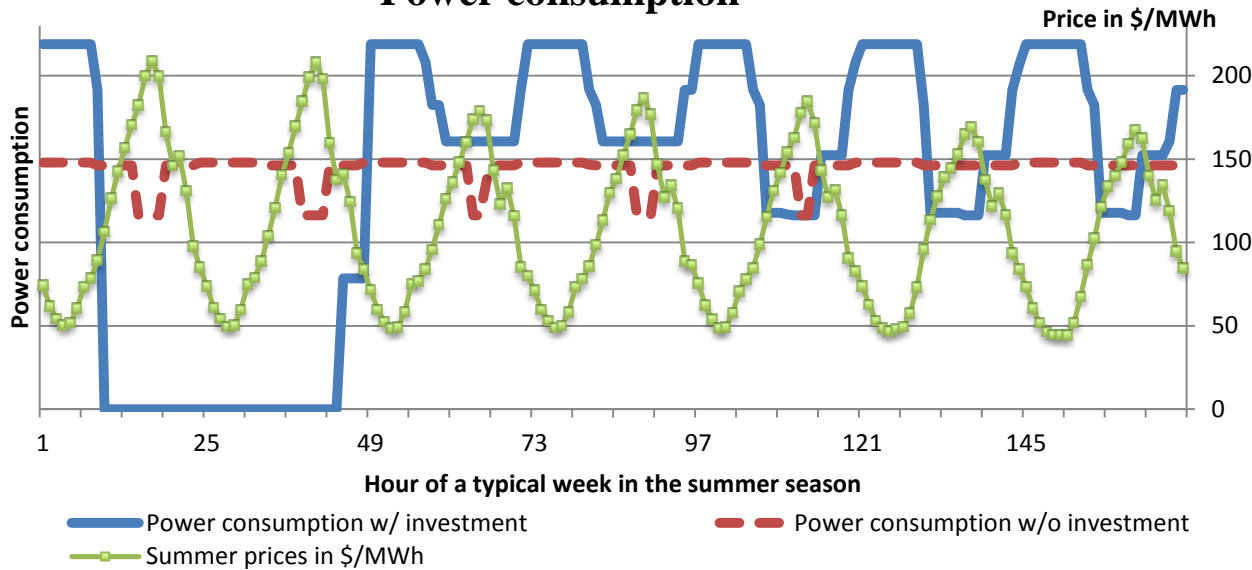


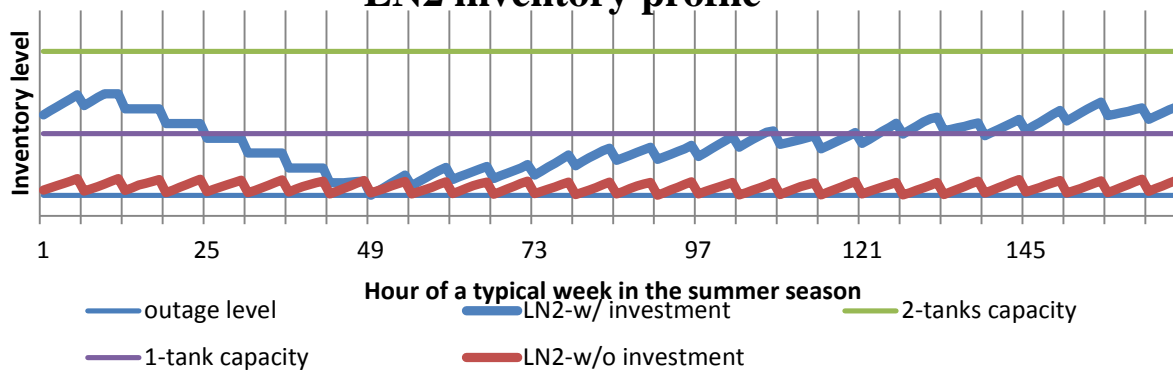
Fig. 2: State graph superstructure

- The resulting MILP has **191,861 constraints** and **161,293 variables (18,826 binary.)**
- Solution time: **38.5 minutes** (GAMS 23.6.2, GUROBI 4.0.0, Intel i7 (2.93GHz) with 4GB RAM)

Power consumption



LN2 inventory profile



Remarks on case study

- **Annualized costs: \$5,700k/yr**
- **Annualized savings: \$400k/yr**
- Buy **new liquefier** in the first time period (annualized investment costs: \$300k/a)
- Buy **additional LN2 storage tank** (\$25k/a)
- **Don't upgrade** existing equipment (\$200k/a)
- Take-away message on operational level: *Reduce production when prices are high and build up LN2 when prices are low.*
- Utilization of existing equipment: 97%.

- **The uncertainty challenge:**

*Short term uncertainties: **robust optimization***

Computation time comparable to deterministic models

*Long term uncertainties: **stochastic programming***

Computation time one to two orders of magnitude larger than deterministic models

Global Sourcing Project with Uncertainties

You, Wassick, Grossmann (2009)

- Given
 - ◆ Initial inventory
 - ◆ Inventory holding cost and throughput cost
 - ◆ **Transport times** of all the transport links
 - ◆ **Uncertain production reliability** and **demands**
- Determine
 - ◆ Inventory levels, transportation and sale amounts



~ 100 facilities
~ 1,000 customers
~ 25,000 shipping links/modes

- **Objective: Minimize Cost**

Two-stage stochastic MILP model
1000 scenarios (Monte Carlo sampling)



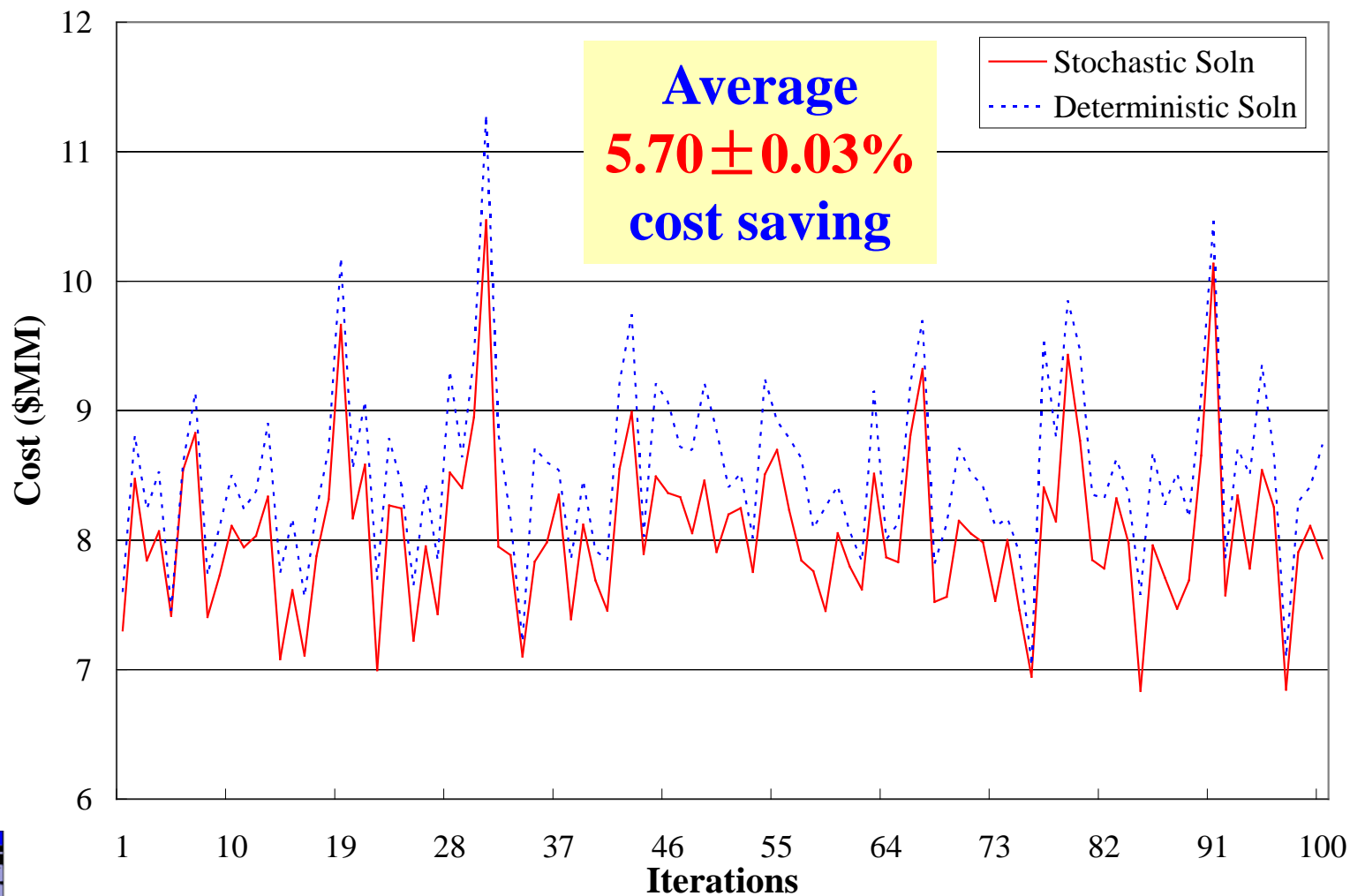
MILP Problem Size

Case Study 1	Deterministic Model
# of Constraints	62,187
# of Cont. Var.	89,014
# of Disc. Var.	7

- ◆ Impossible to solve directly
- ◆ takes 5 days by using standard L-shaped Benders
- ◆ only 20 hours with multi-cut version Benders
- ◆ 30 min if using 50 parallel CPUs and multi-cut version



Stochastic Planner vs Deterministic Planner



➤ Offshore oilfield having several reservoirs **under uncertainty**

Tarhan, Grossmann (2010)

➤ **Maximize the expected net present value (ENPV)** of the project

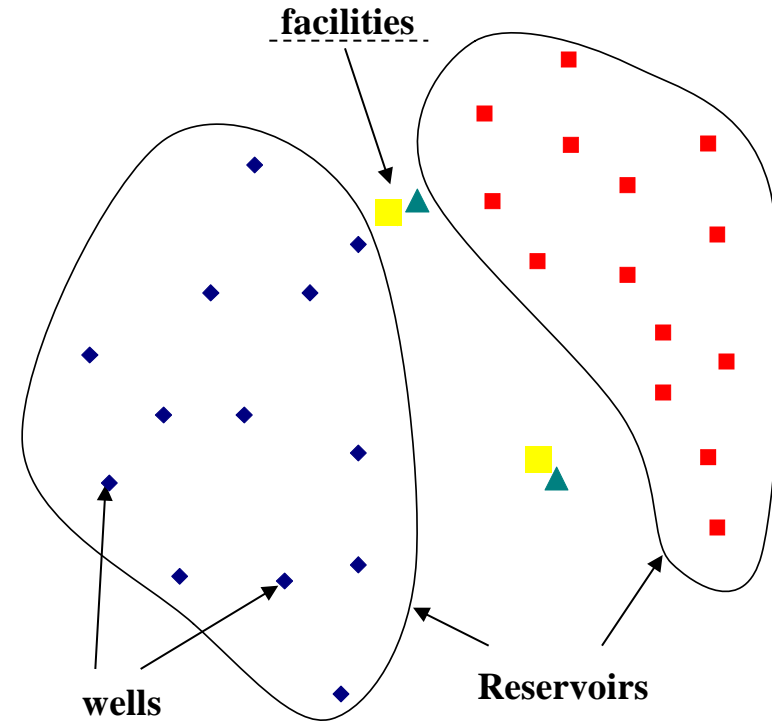
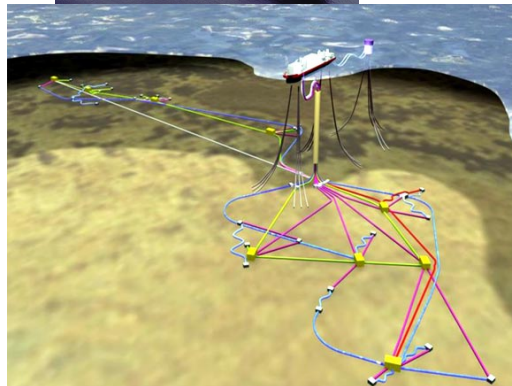
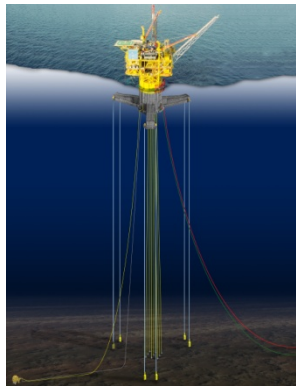
Decisions:

- Number and capacity of TLP/FPSO facilities
- Installation schedule for facilities
- Number of sub-sea/TLP wells to drill
- Oil production profile over time

TLP



FPSO



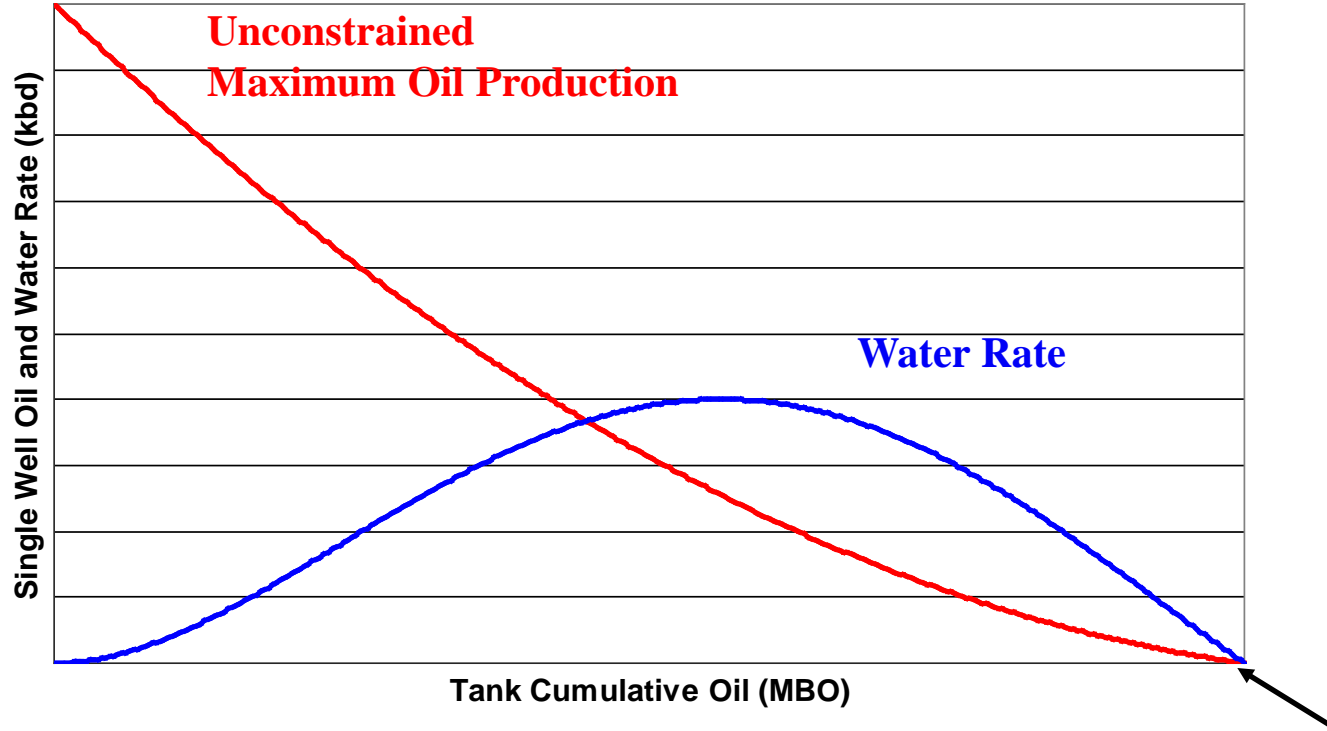
Uncertainty:

- **Initial productivity per well**
- **Size of reservoirs**
- **Water breakthrough time for reservoirs**

Non-linear Reservoir Model

Initial oil
production

Assumption: All wells in the same reservoir are identical.



Size of the reservoir

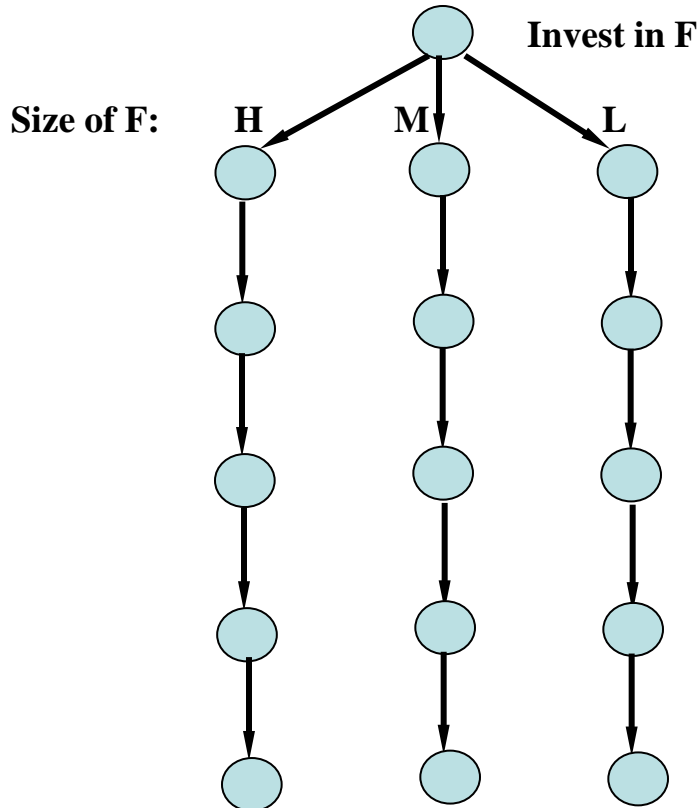
Uncertainty is represented by discrete distributions functions

Decision Dependent Scenario Trees

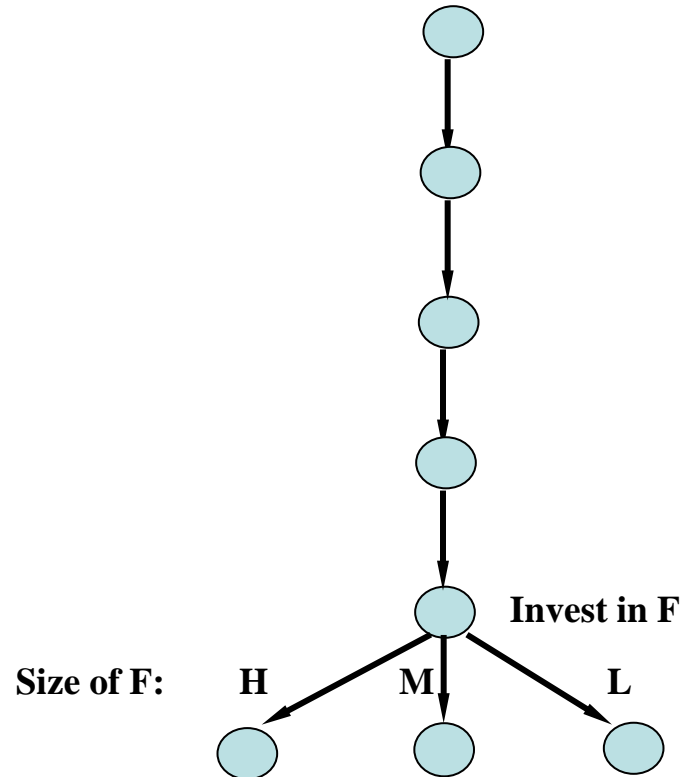
(Endogeneous uncertainties)

Assumption: Uncertainty in a field resolved as soon as WP installed at field

Invest in F in year 1

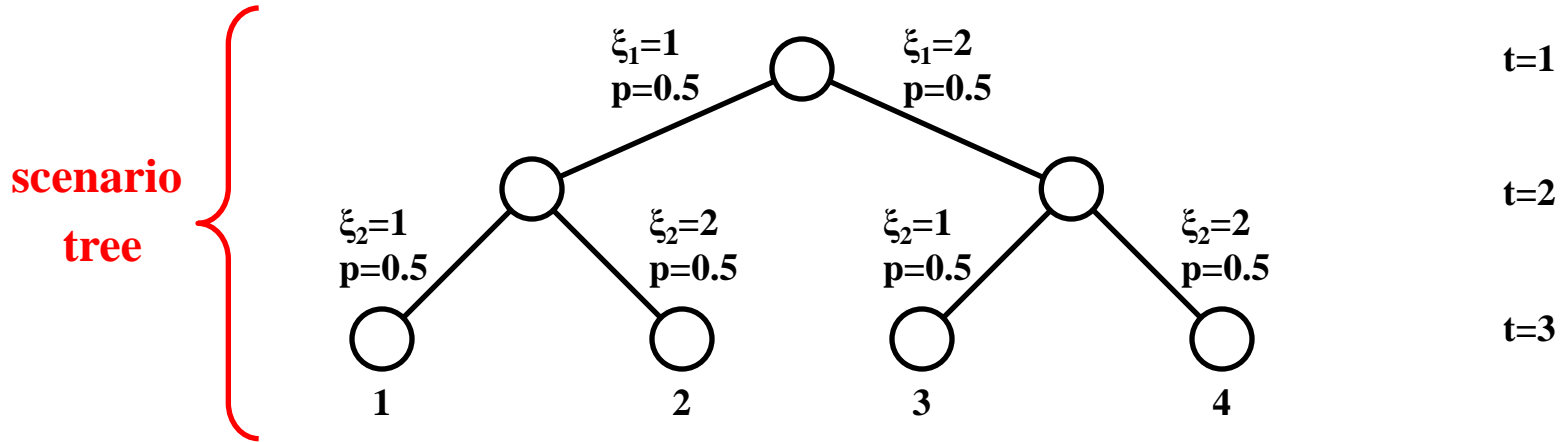


Invest in F in year 5

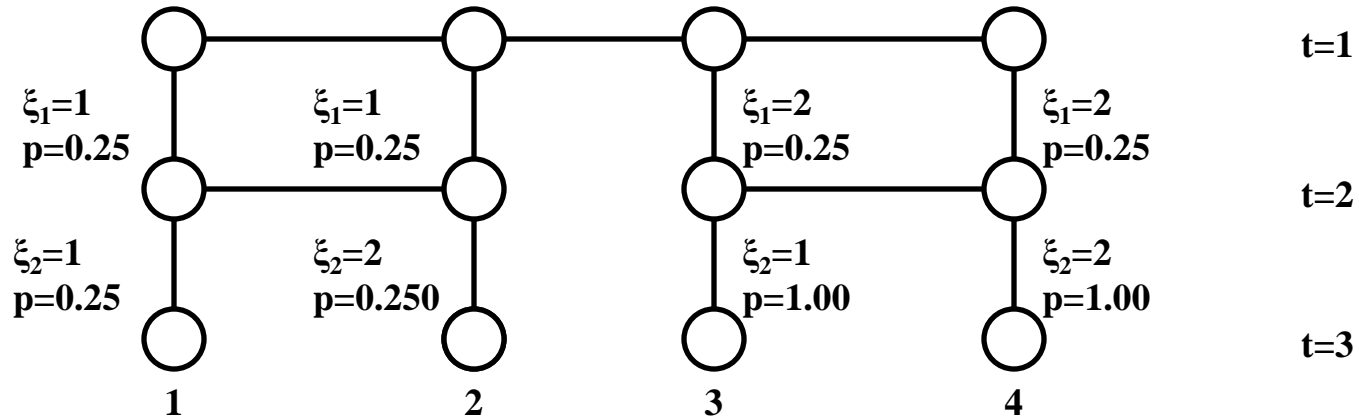


Scenario tree

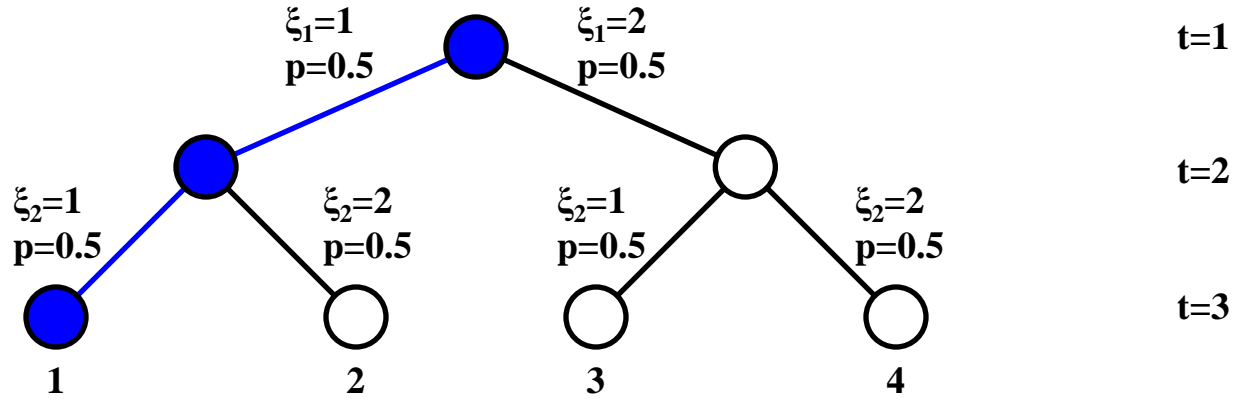
Not unique: Depends on timing of investment at uncertain fields
 Central to defining a Stochastic Programming Model



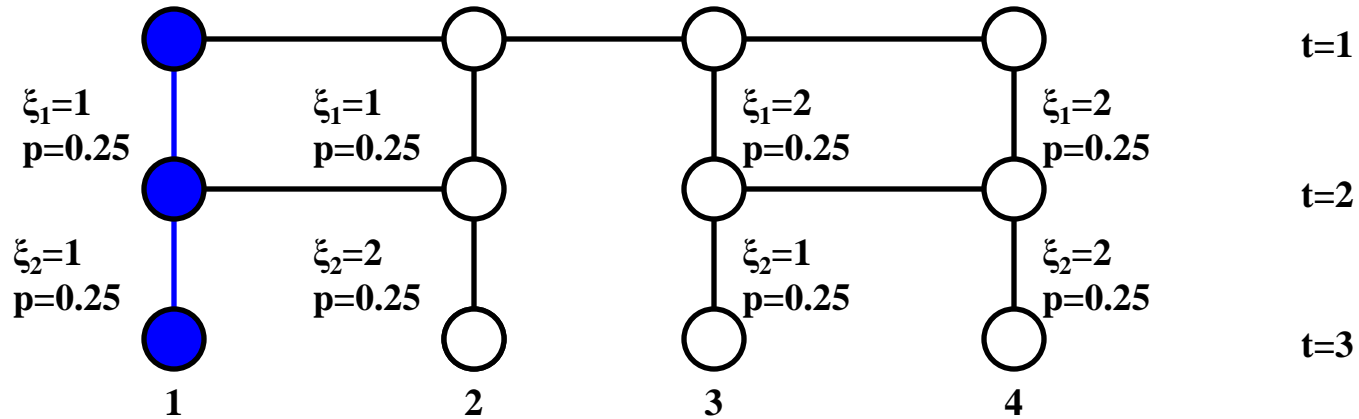
Alternative and equivalent scenario tree structure (Ruszczynski, 1997):



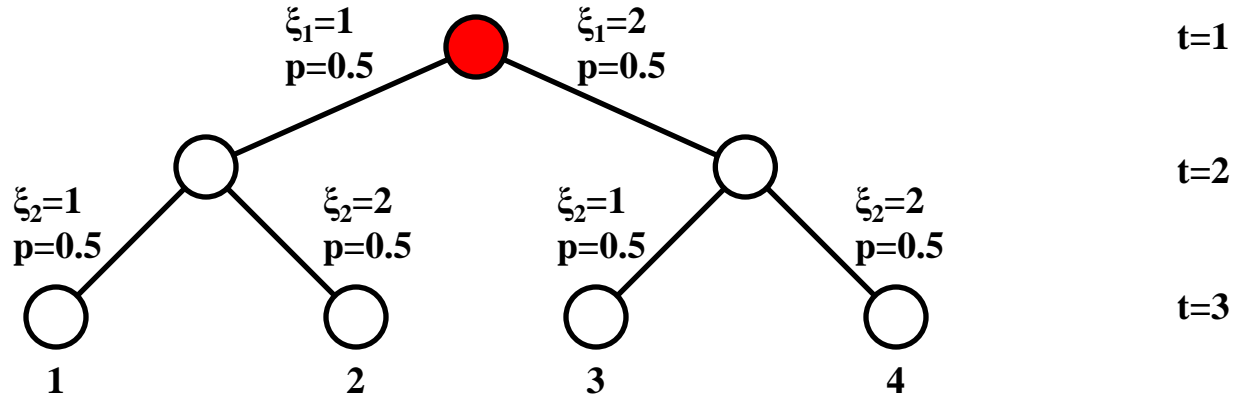
Stochastic Programming



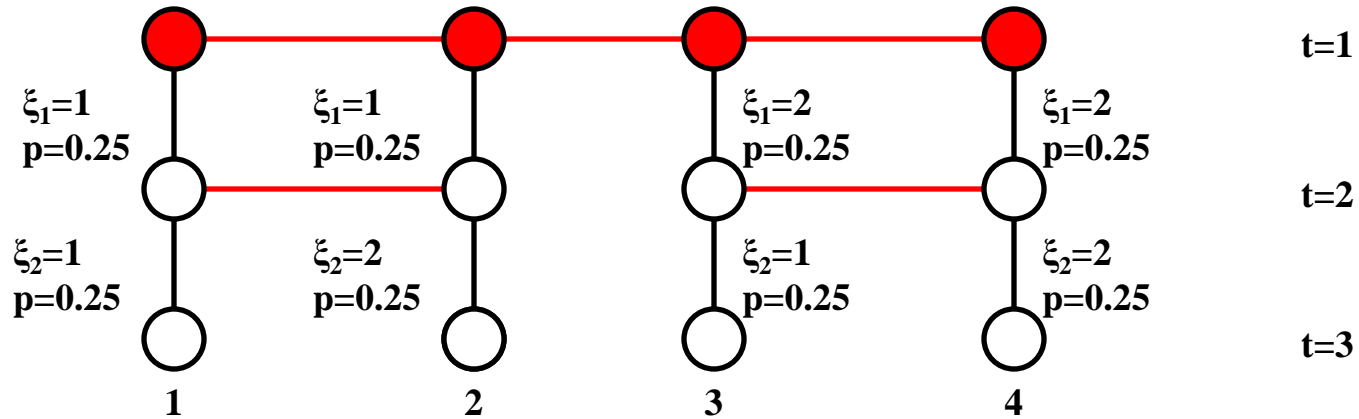
Each scenario is represented by a set of unique nodes



Stochastic Programming

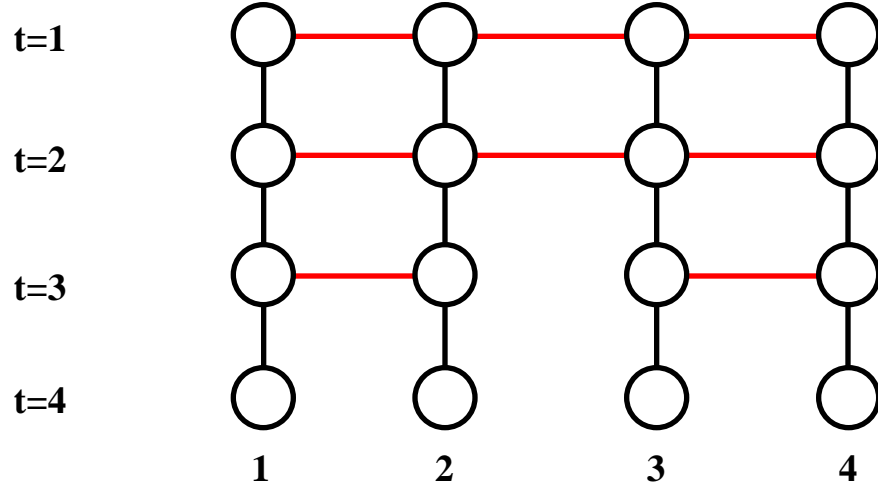
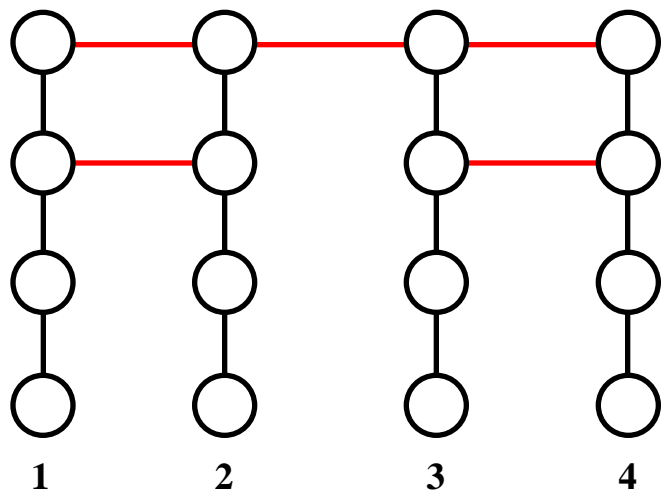
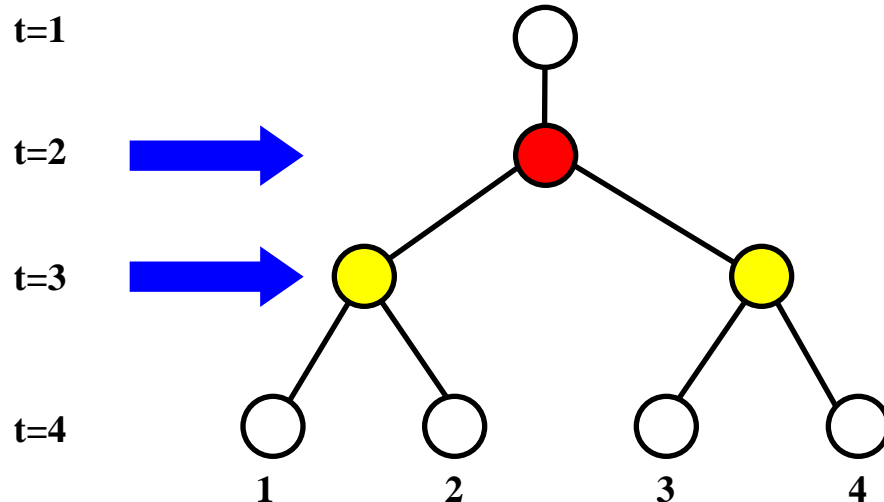
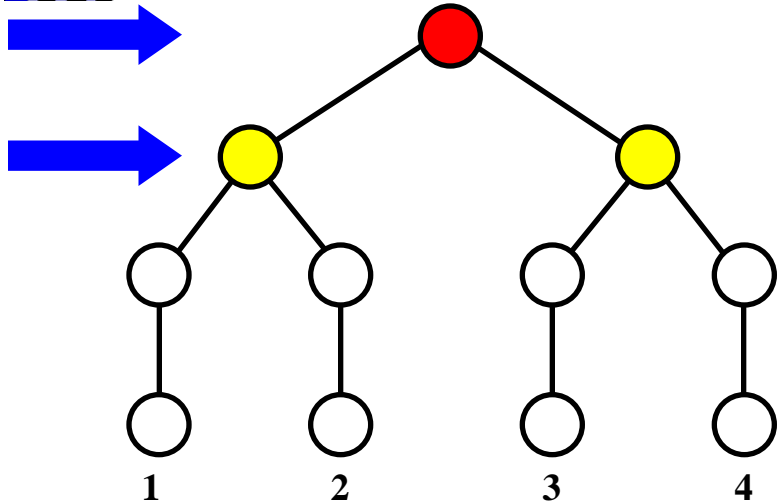


Nodes have **same amount of information** \equiv Nodes are **indistinguishable**



Non-anticipativity constraints

Representation of Decision-Dependence Using Scenario Tree



Maximize.. Probability weighted average of NPV over uncertainty scenarios
subject to

- Equations about economics of the model
 - Surface constraints
 - **Non-linear equations related to reservoir performance**
 - Logic constraints relating decisions
- if there is a TLP available, a TLP well can be drilled

Every
scenario,
time period

- Non-anticipativity constraints

Non-anticipativity prevents a decision being taken now from using information that will only become available in the future

Every pair
scenarios,
time period

Disjunctions (conditional constraints)

Problem size MINLP increases exponentially with number of time periods and scenarios



Decomposition algorithm:
Lagrangean relaxation & Branch and Bound

Formulation of Lagrangean dual

Relaxation

- Relax disjunctions, logic constraints
- Penalty for equality constraints

$b_{uf}^{s,s'}, y_{\lambda}^{s,s'}, d_{\lambda}^{s,s'}$:

Lagrange Multipliers

$$\text{Max } \sum_s p^s \left[\sum_t \left(c_{1t} q_t^s + c_{2t} d_t^s + c_{3t} y_t^s + \sum_{uf} c_{4t,uf} b_{uf,t}^s \right) \right] + \sum_{(s,s')} \left[\sum_{uf} b_{\lambda_{uf}}^{s,s'} (b_{uf,1}^s - b_{uf,1}^{s'}) + y_{\lambda}^{s,s'} (y_1^s - y_1^{s'}) + d_{\lambda}^{s,s'} (d_1^s - d_1^{s'}) \right]$$

$$\sum_{\tau=1}^t \left(A_{\tau}^s q_{\tau}^s + B_{\tau}^s d_{\tau}^s + C_{\tau}^s y_{\tau}^s + \sum_{uf} D_{uf,\tau}^s b_{uf,\tau}^s \right) \leq a_t^s \quad \forall(t, s)$$

$$\begin{bmatrix} q_t^s & = & q_t^{s'} \\ d_{t+1}^s & = & d_{t+1}^{s'} \\ y_{t+1}^s & = & y_{t+1}^{s'} \\ b_{uf,t+1}^s & = & b_{uf,t+1}^{s'} \quad \forall uf \end{bmatrix} \quad \forall \left[\neg Z_t^{s,s'} \right] \quad \forall(t, s, s')$$

$$Z_t^{s,s'} \Leftrightarrow \bigwedge_{uf \in \mathcal{D}(s,s')} \left[\bigwedge_{\tau=1}^t (\neg b_{uf,\tau}^s) \right] \quad \forall(t, s, s')$$

$$\begin{aligned} b_{uf,1}^s &= b_{uf,1}^{s'} && \forall(uf, s, s') \\ d_1^s &= d_1^{s'} && \forall(s, s') \\ y_1^s &= y_1^{s'} && \forall(s, s') \end{aligned}$$

Non-anticipativity constraints

One Reservoir Example

Optimize the planning decisions for an oilfield having **single reservoir** for **10 years**.

Decisions:

Number, capacity and installation schedule of FPSO/TLP facilities

Number and drilling schedule of sub-sea/TLP wells

Oil production profile over time

Uncertain Parameters (Discrete Values)	Scenarios							
	1	2	3	4	5	6	7	8
Initial Productivity <u>per</u> well (kbd)	10	10	20	20	10	10	20	20
Reservoir Size (Mbbbl)	300	300	300	300	1500	1500	1500	1500
Water Breakthrough Time Parameter	5	2	5	2	5	2	5	2

Construction Lead Time (years)	Wells		Facilities		
	TLP	Sub-sea	TLP	Small FPSO	Large FPSO
	1	1	1	2	4

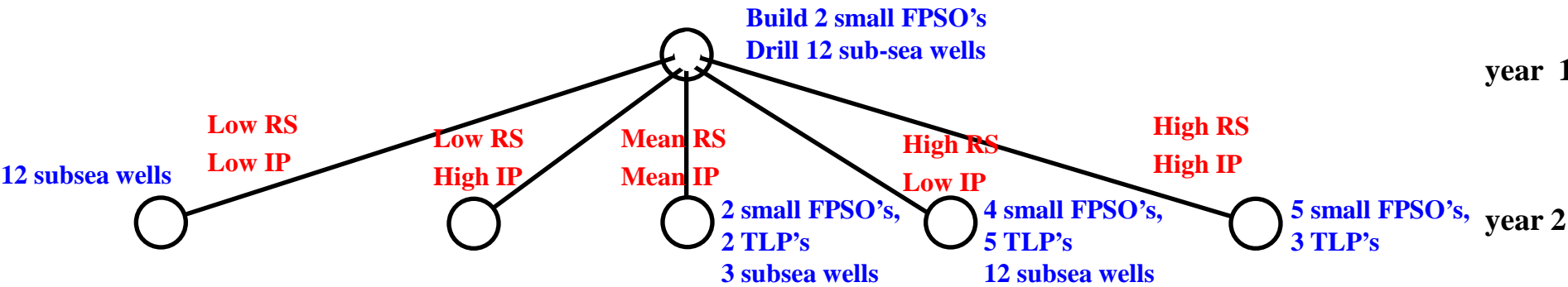


Multistage Stochastic Programming Approach



RS: Reservoir size
IP: Initial Productivity
BP: Breakthrough Parameter

$$E[NPV] = \$4.92 \times 10^9$$

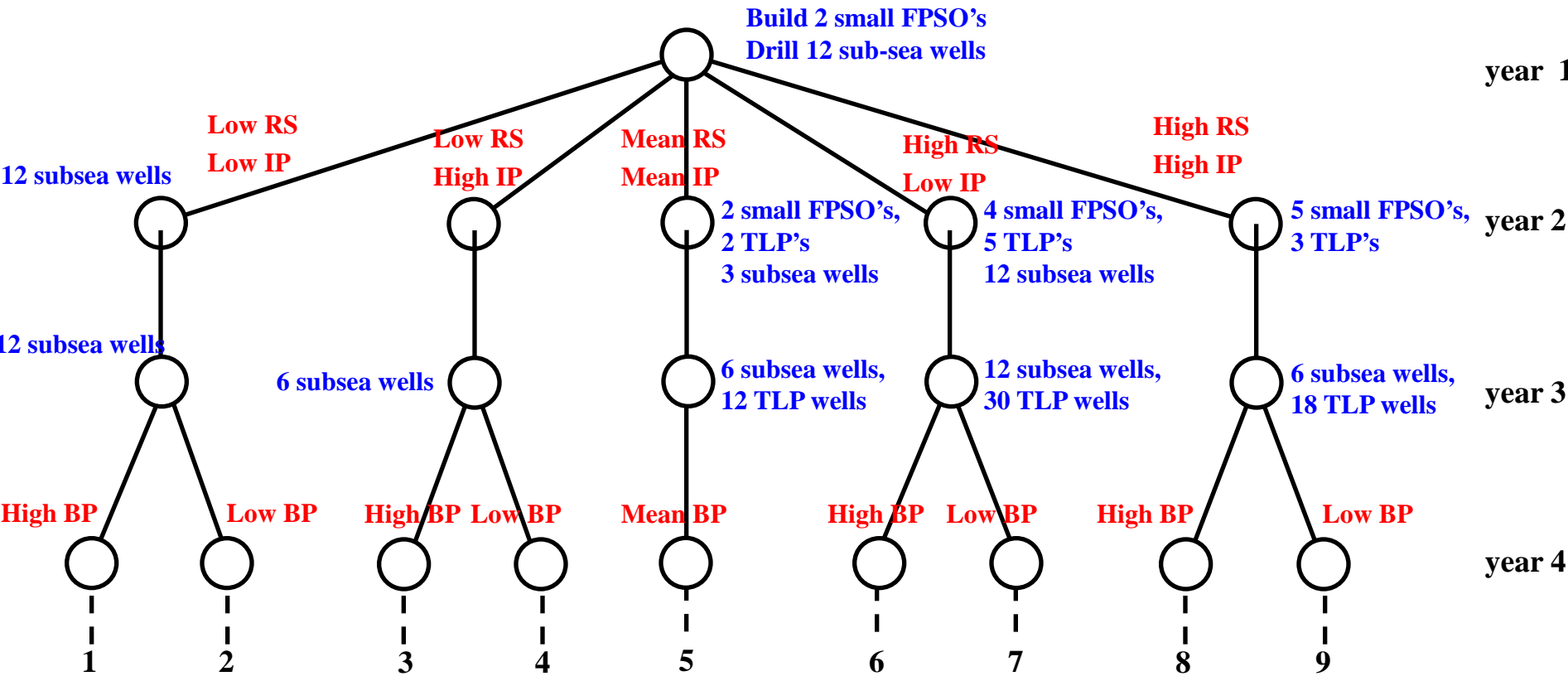


Solution proposes building **2** small FPSO's in the first year and then add new facilities / drill wells (**recourse action**) depending on the positive or negative outcomes.

Multistage Stochastic Programming Approach

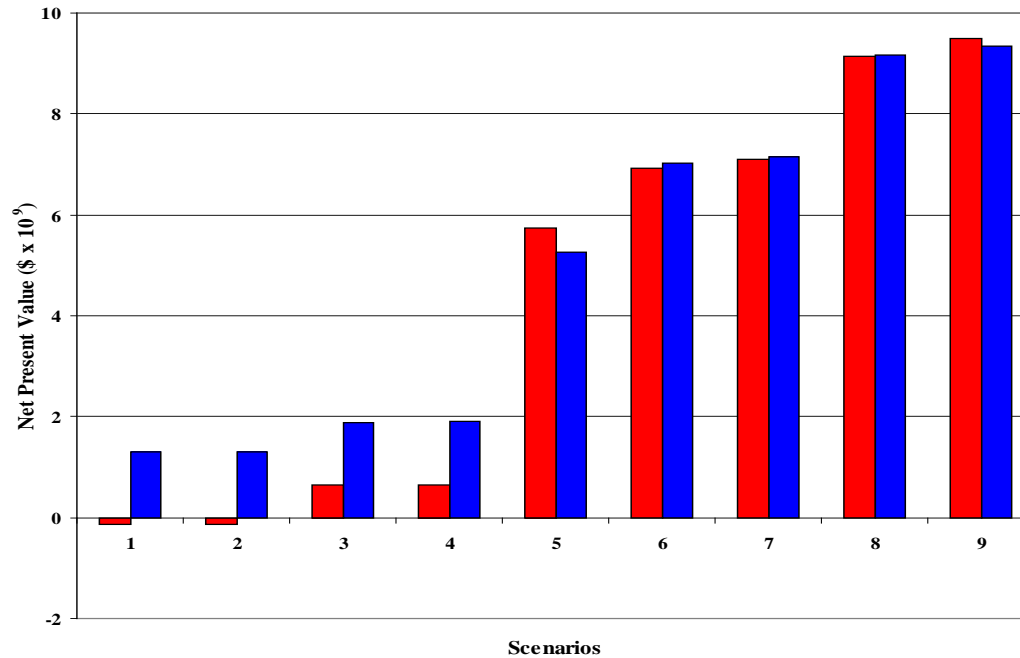
RS: Reservoir size
IP: Initial Productivity
BP: Breakthrough Parameter

$E[NPV] = \$4.92 \times 10^9$



Solution proposes building **2** small FPSO's in the first year and then add new facilities / drill wells (**recourse action**) depending on the positive or negative outcomes.

Distribution of Net Present Value



Deterministic Mean Value = \$4.38 x 10⁹



Multistage Stoch Progr = \$4.92 x 10⁹ => 12% higher and more robust

Computation: Algorithm 1: 120 hrs; Algorithm 2: 5.2 hrs

Nonconvex MINLP: 1400 discrete vars, 970 cont vars, 8090 Constraints

Economics vs. performance?

Multiobjective Optimization Approach

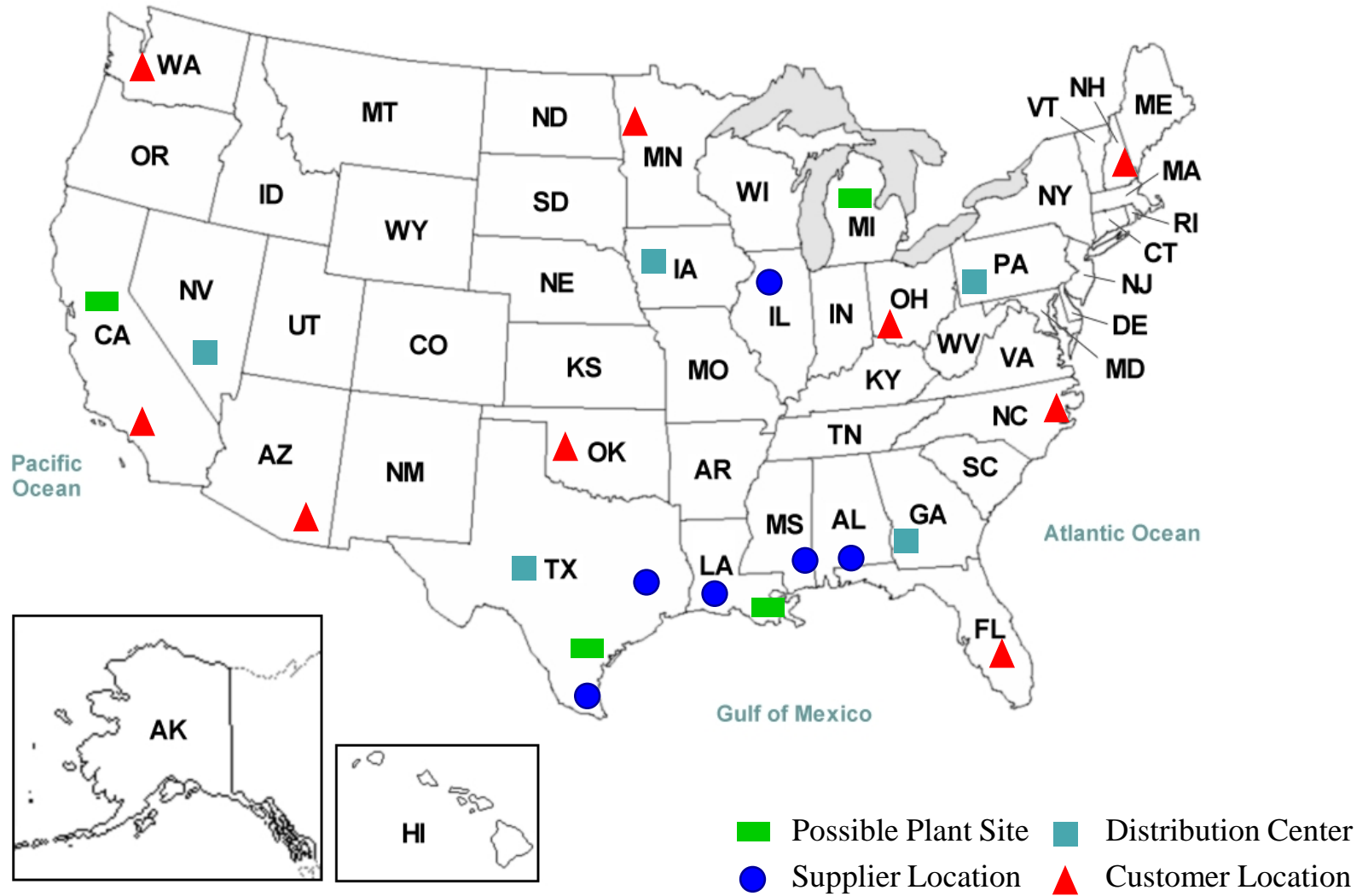
Economics vs Environmental:

Guillen-Gozalbez, Grossmann (2010)

Pinto-Varela, Barbosa-Póvoa and A.Q. Novais (2011)

Objective: design supply chain polystyrene resins under **responsive** and **economic** criteria

You, Grossmann (2008)



Production Network of Polystyrene Resins

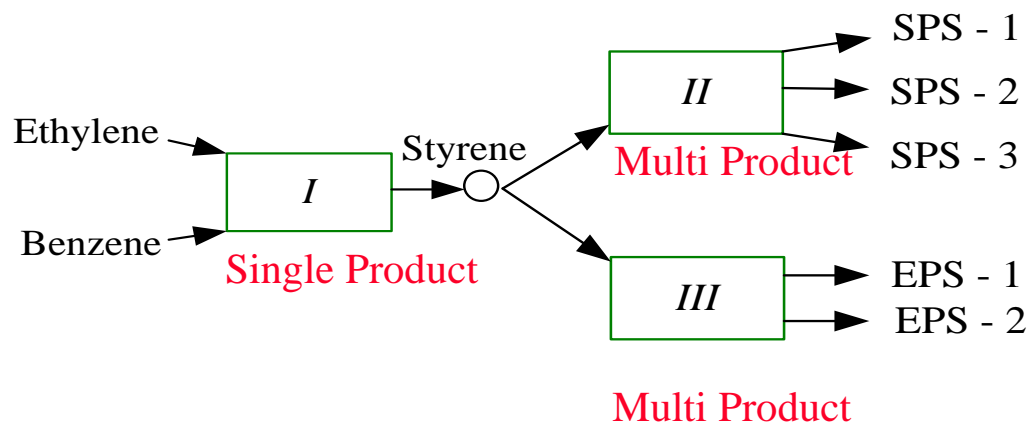
Three types of plants:

Plant I: *Ethylene + Benzene* \longrightarrow *Styrene* (1 products)

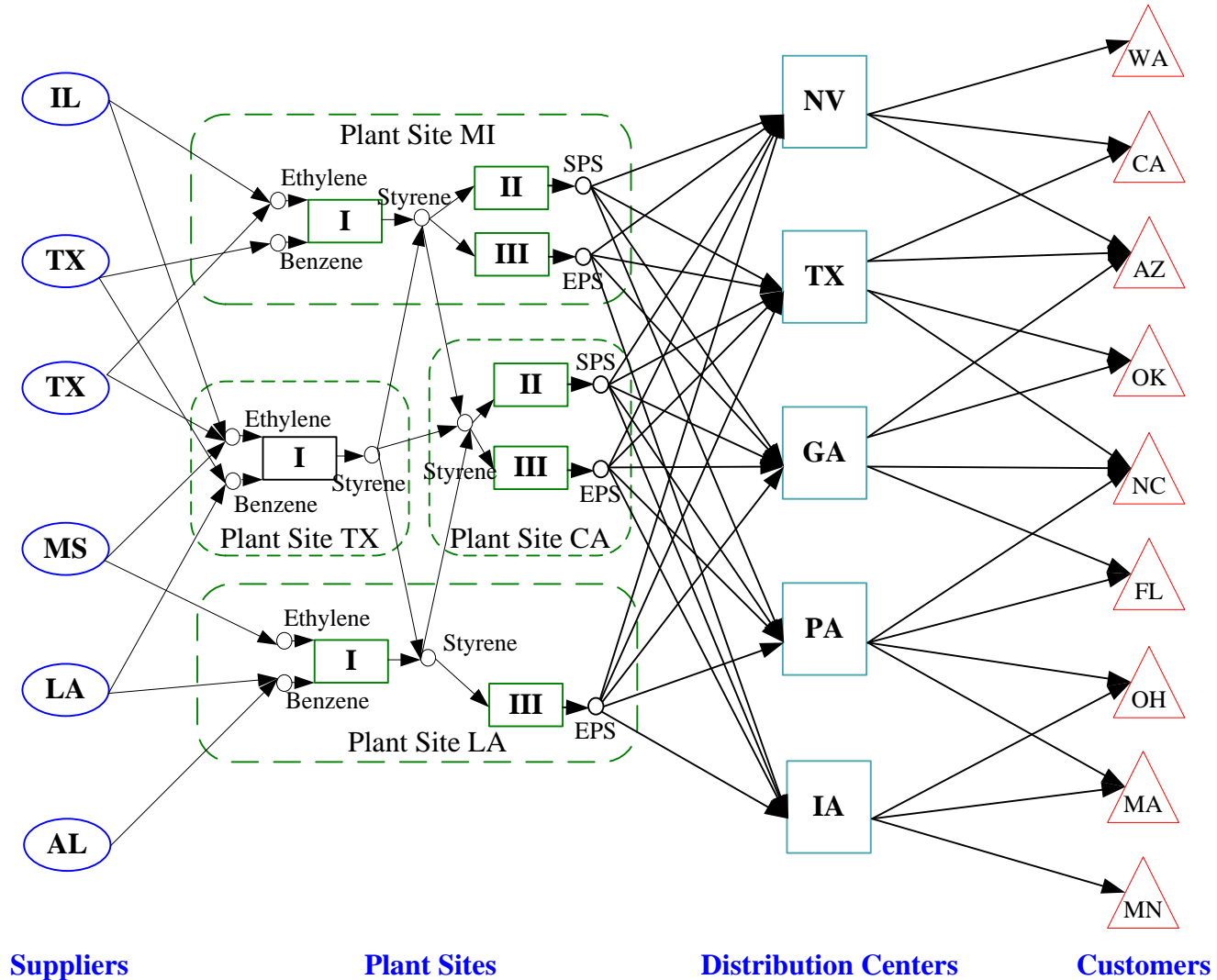
Plant II: *Styrene* \longrightarrow *Solid Polystyrene (SPS)* (3 products)

Plant III: *Styrene* \longrightarrow *Expandable Polystyrene (EPS)* (2 products)

Basic Production Network



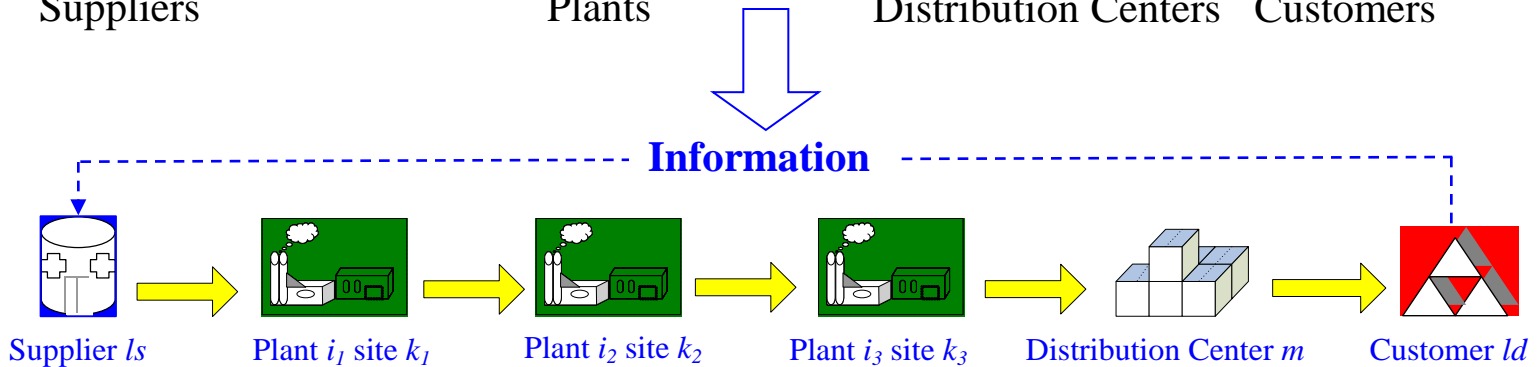
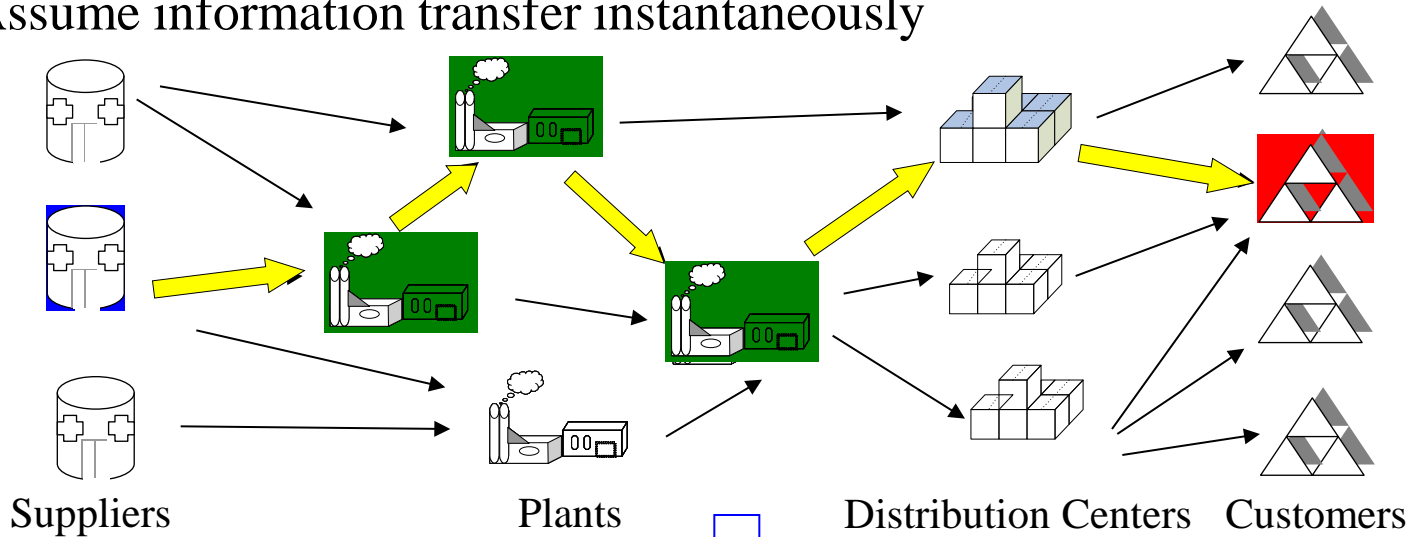
Potential Network Superstructure



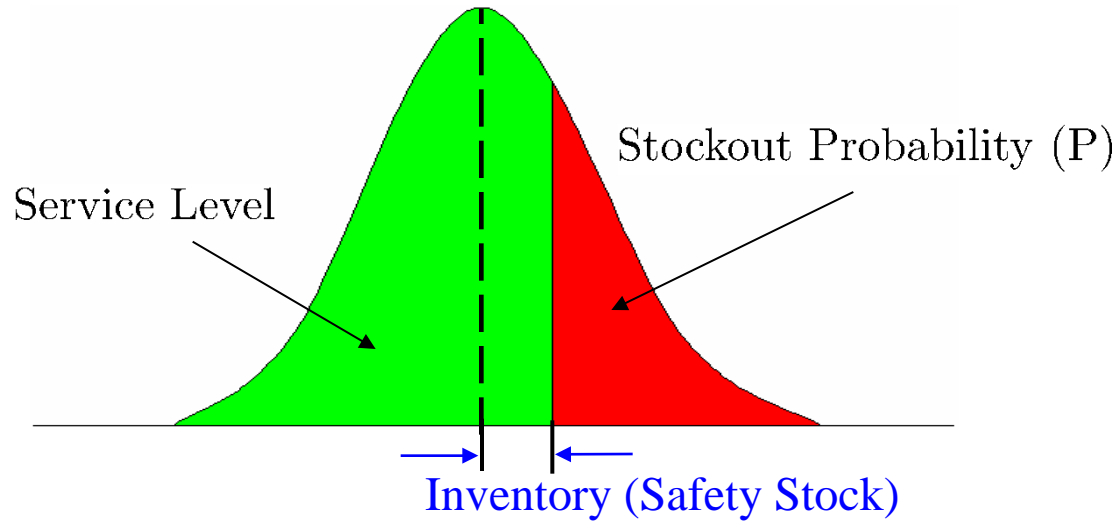
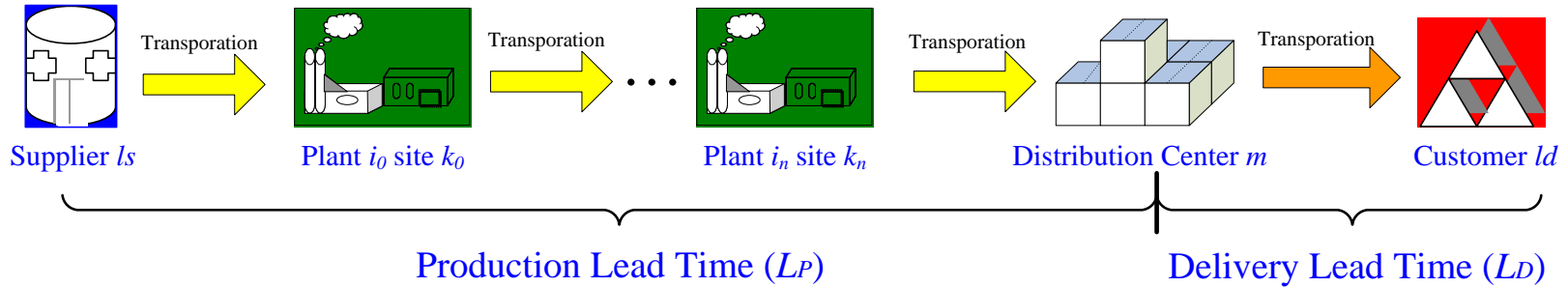
Responsiveness - Lead Time

Lead Time for A Linear Supply Chain

- A supply chain network = \sum Linear supply chains
- ♦ Assume information transfer instantaneously



Lead Time under Demand Uncertainty



$$\text{Expected Lead Time} = L_D + P(\text{Stockout}) \cdot L_P$$

Bi-criterion Multiperiod MINLP Formulation

Choose Discrete (0-1), continuous variables

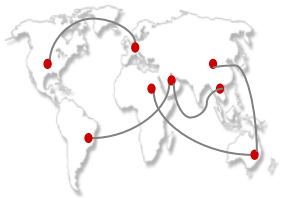
- Objective Function:

- ◆ Max: Net Present Value
 - ◆ Min: Expected Lead time
- } Bi-criterion

- Constraints:

- ◆ Network structure constraints

- Suppliers – plant sites Relationship
- Plant sites – Distribution Center
- Input and output relationship of a plant
- Distribution Center – Customers
- Cost constraint



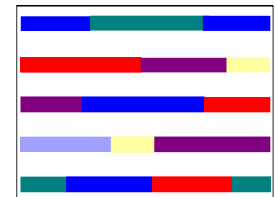
- ◆ Operation planning constraints

- Production constraint
- Capacity constraint
- Mass balance constraint
- Demand constraint
- Upper bound constraint



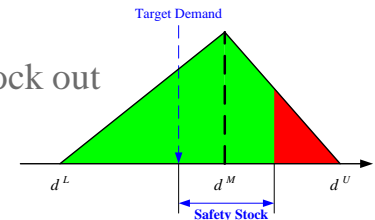
- ◆ Cyclic scheduling constraints

- Assignment constraint
- Sequence constraint
- Demand constraint
- Production constraint
- Cost constraint

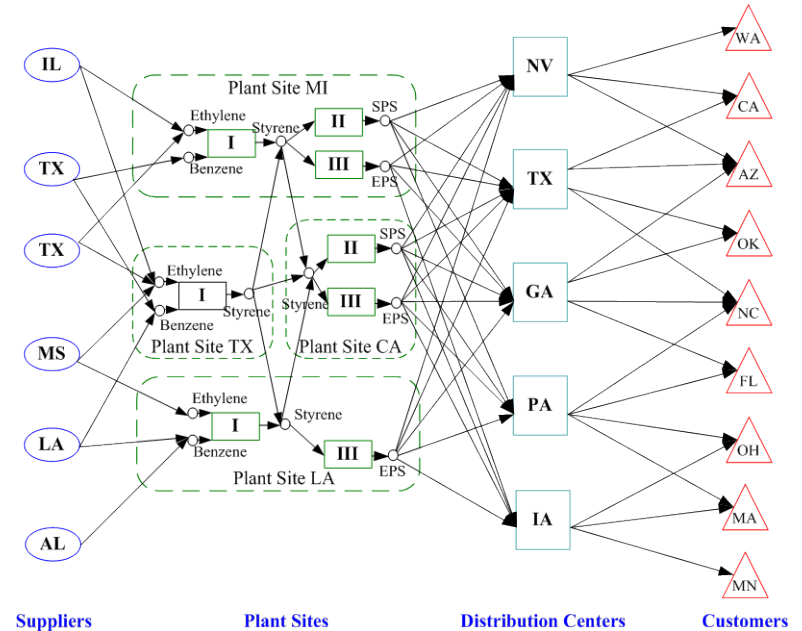
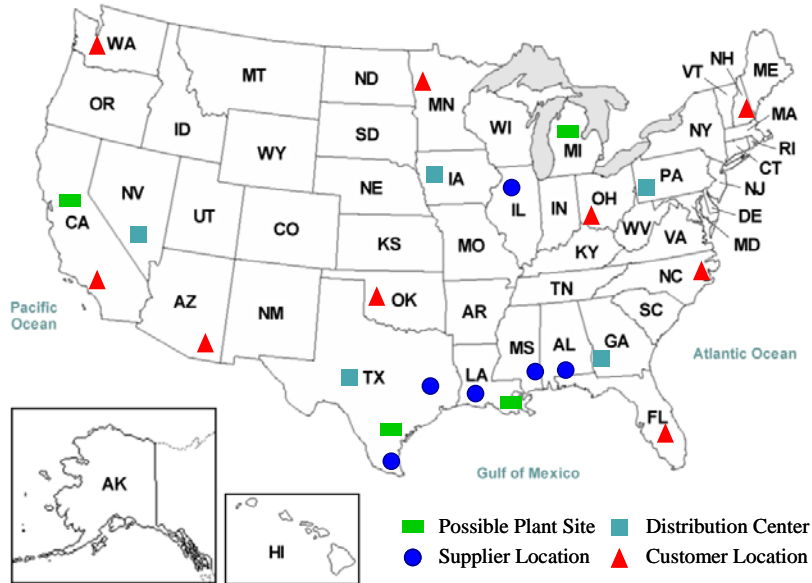


- ◆ Probabilistic constraints

- Chance constraint for stock out (reformulations)



Case Study



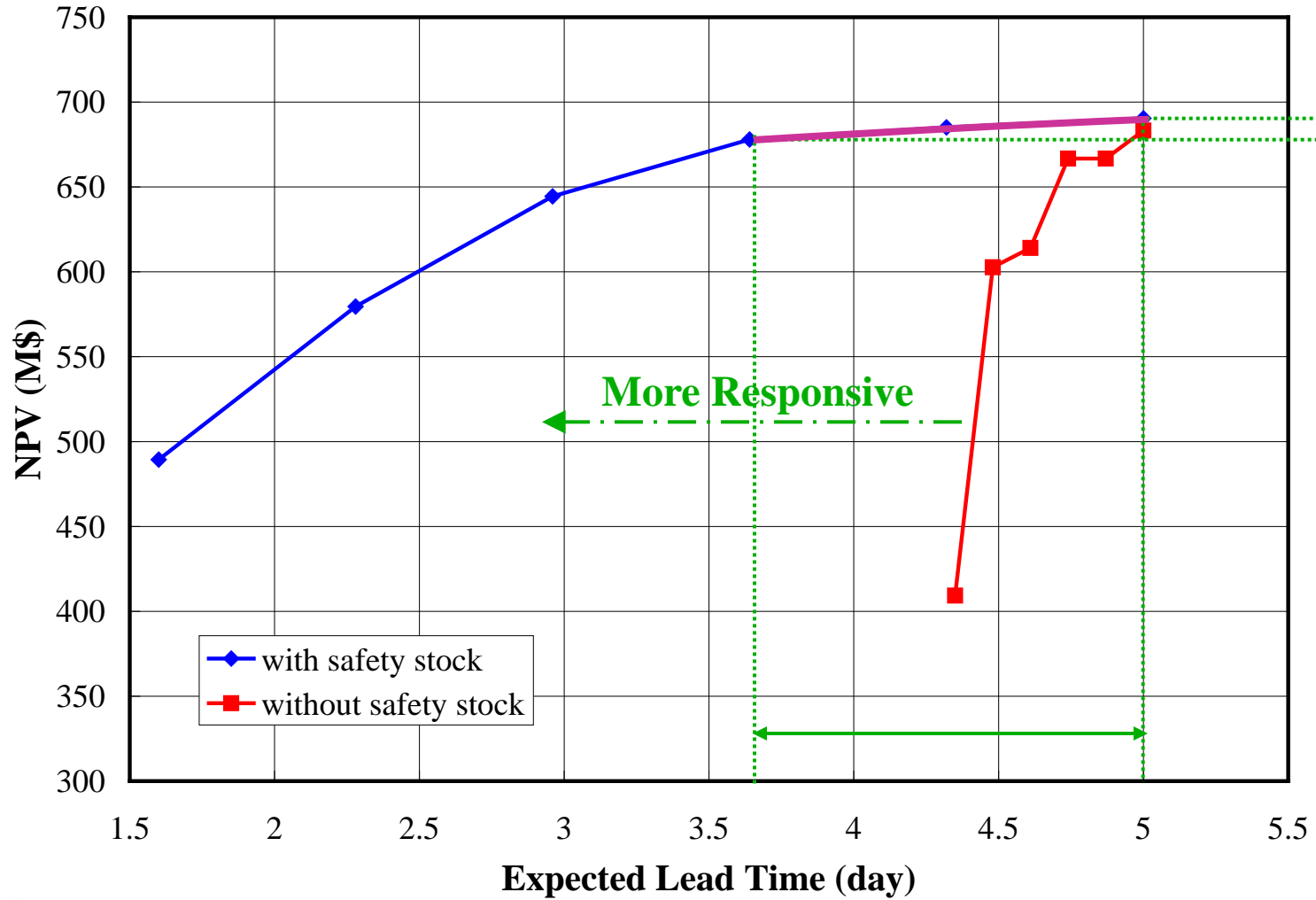
- Problem Size:

- ◆ # of Discrete Variables: 215
- ◆ # of Continuous Variables: 8126
- ◆ # of Constraints: 14617

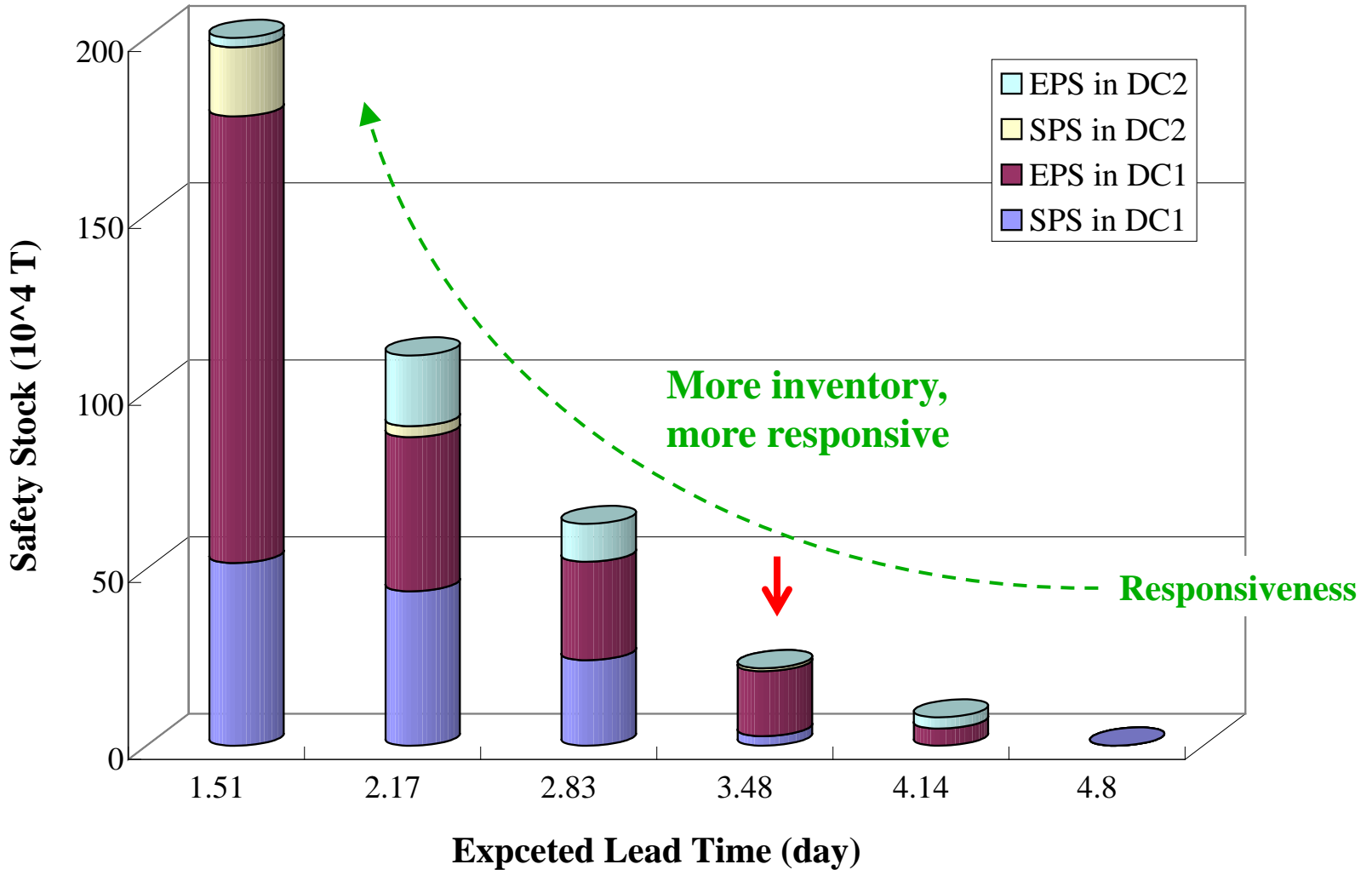
- Solution Time:

- ◆ Solver: GAMS/BARON
- ◆ Direct Solution: > 2 weeks
- ◆ Proposed Algorithm: ~ 4 hours

Pareto Curves – with and without safety stock



Safety Stock Levels - Expected Lead Time



1. Integration of **control** with planning and scheduling

Bhatia, Biegler (1996), Perea, Ydstie, Grossmann (2003), Flores, Grossmann (2006), Prata, Oldenburg, Kroll, Marquardt (2008), Harjunkoski, Nystrom, Horch (2009)

Challenge: Effective solution of Mixed-Integer Dynamic Optimization (MIDO)

2. Optimization of entire supply chains

Challenges:

- **Combining different models (eg maritime and vehicle transportation, pipelines)**
Cafaro, Cerda (2004), Relvas, Matos, Barbosa-Póvoa, Fialho, Pinheiro (2006)
- **Advanced financial models**
Van den Heever, Grossmann (2000), Guillén, Badell, Espuña, Puigjaner (2006),

3. Design and Operation of Sustainable Supply Chains

Challenges:

Biofuels, Energy, Environmental

Elia, Baliban, Floudas (2011) Guillén-Gosálbez (2011), You, Tao, Graziano, Snyder (2011)

1. Enterprise-wide Optimization area of great industrial interest

Great economic impact for effectively managing complex supply chains

2. Key components: Planning and Scheduling

Modeling challenge:

Multi-scale modeling (temporal and spatial integration)

3. Computational challenges lie in:

a) Large-scale optimization models (decomposition, advanced computing)

b) Handling uncertainty (stochastic programming)